A fast algorithm for simultaneous computation of matching-polynomials of simple graphs

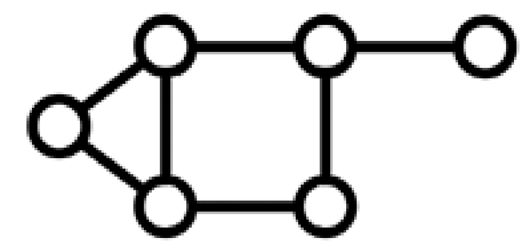
Mario Weitzer

July 1, 2010

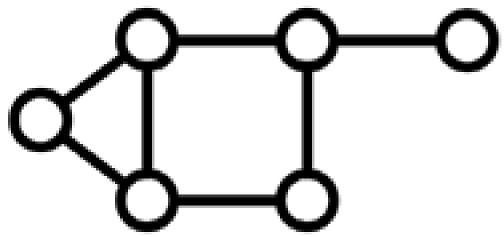
Definitions

Definitions

Simple Graph:
Undirected graph without loops or multiple edges

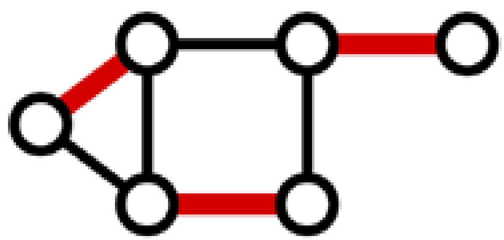


Definitions



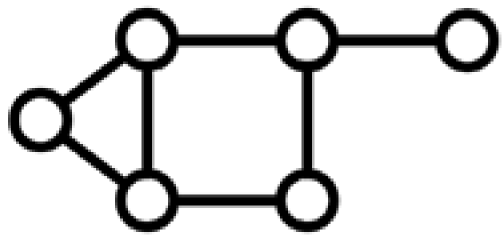
Simple Graph:
Undirected graph without loops or multiple edges

Matching of a graph:
Subset of the edge set such that no two edges share a common vertex (Independet edge set)



k - Matching of a graph:
Matching of cardinality k

Definitions



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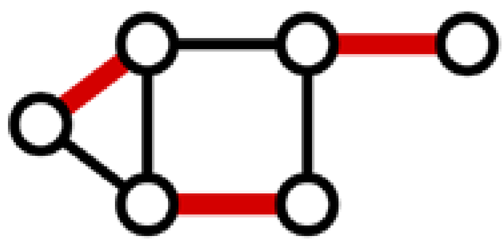
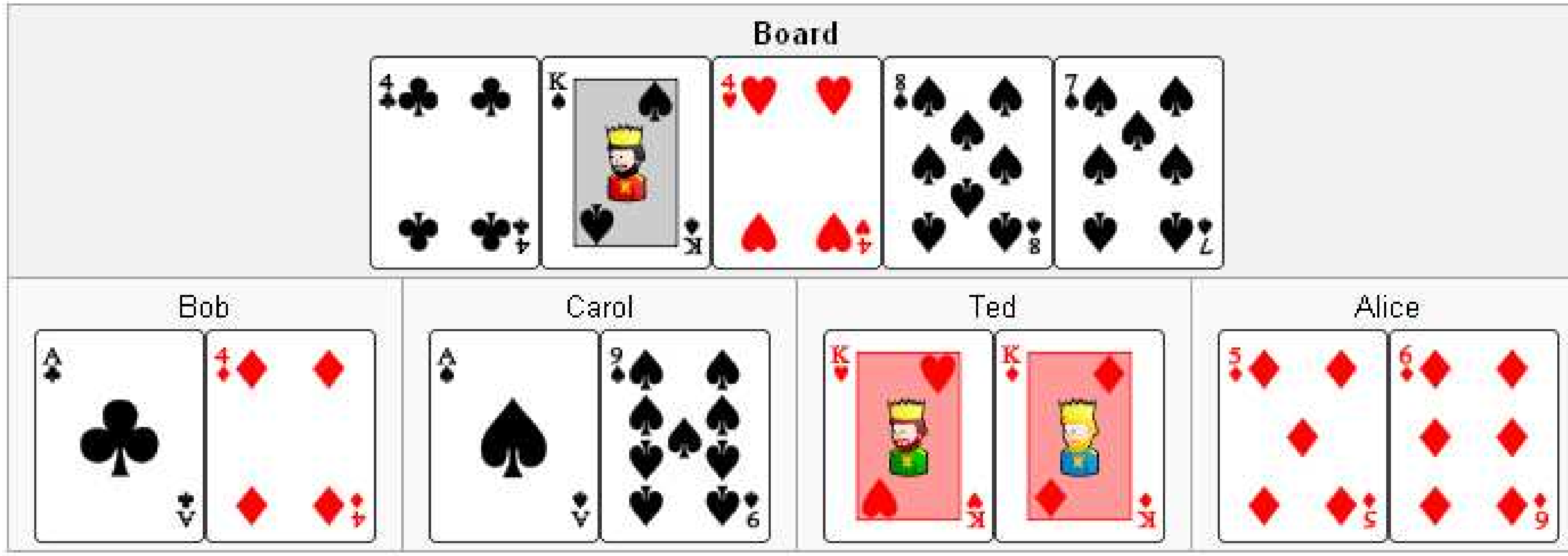


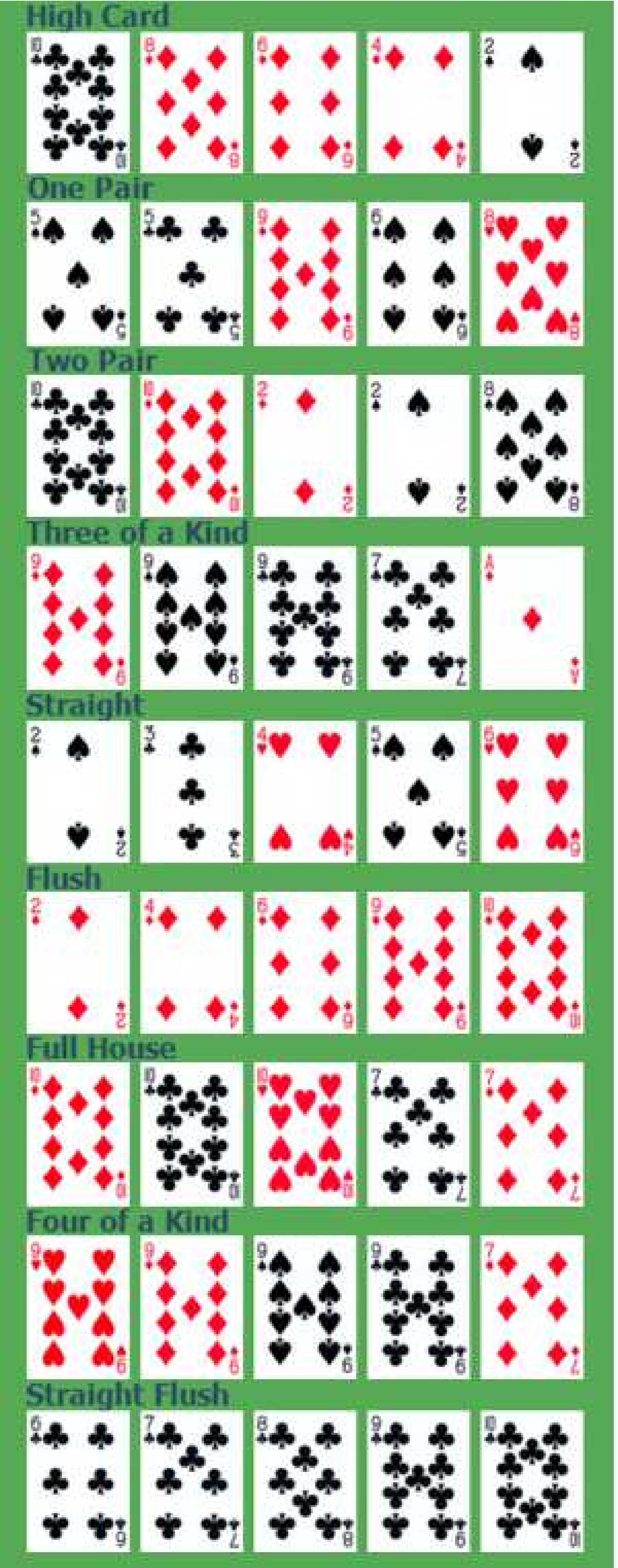
k - Matching of a graph:
Matching of cardinality k

Matching polynomial of a graph:
Coefficient of x^k is number of k - Matchings

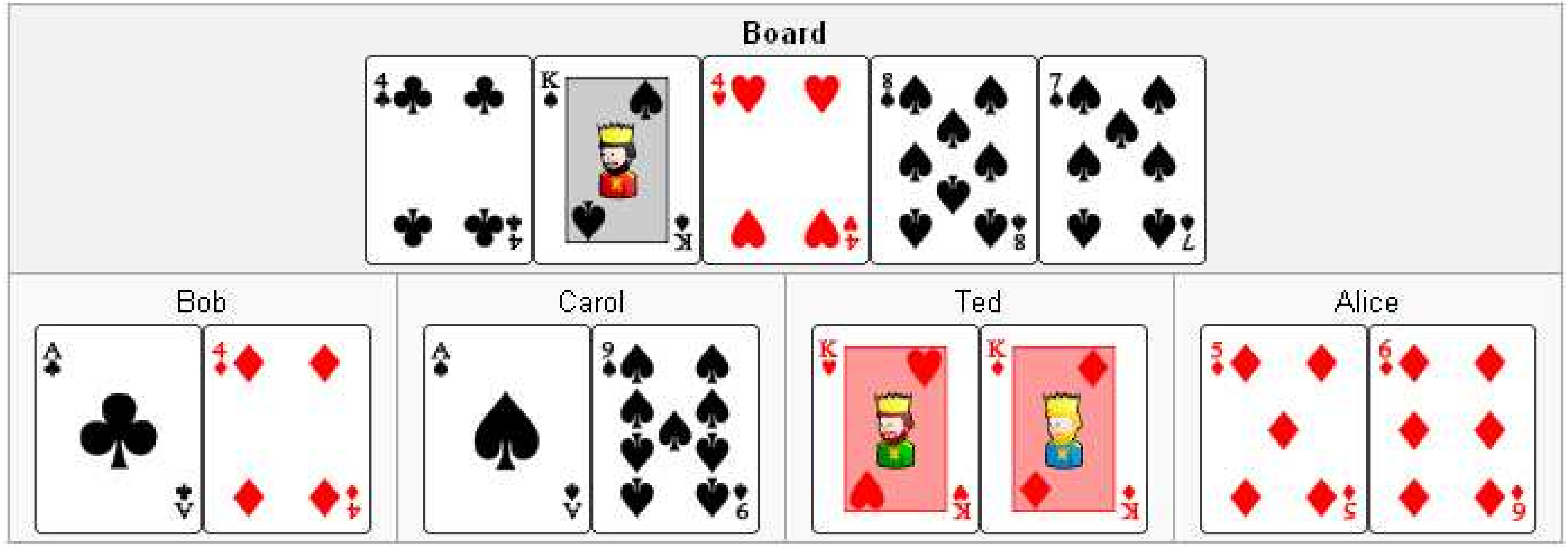
Connection to probabilities in Texas Hold'em Poker?

Sample showdown

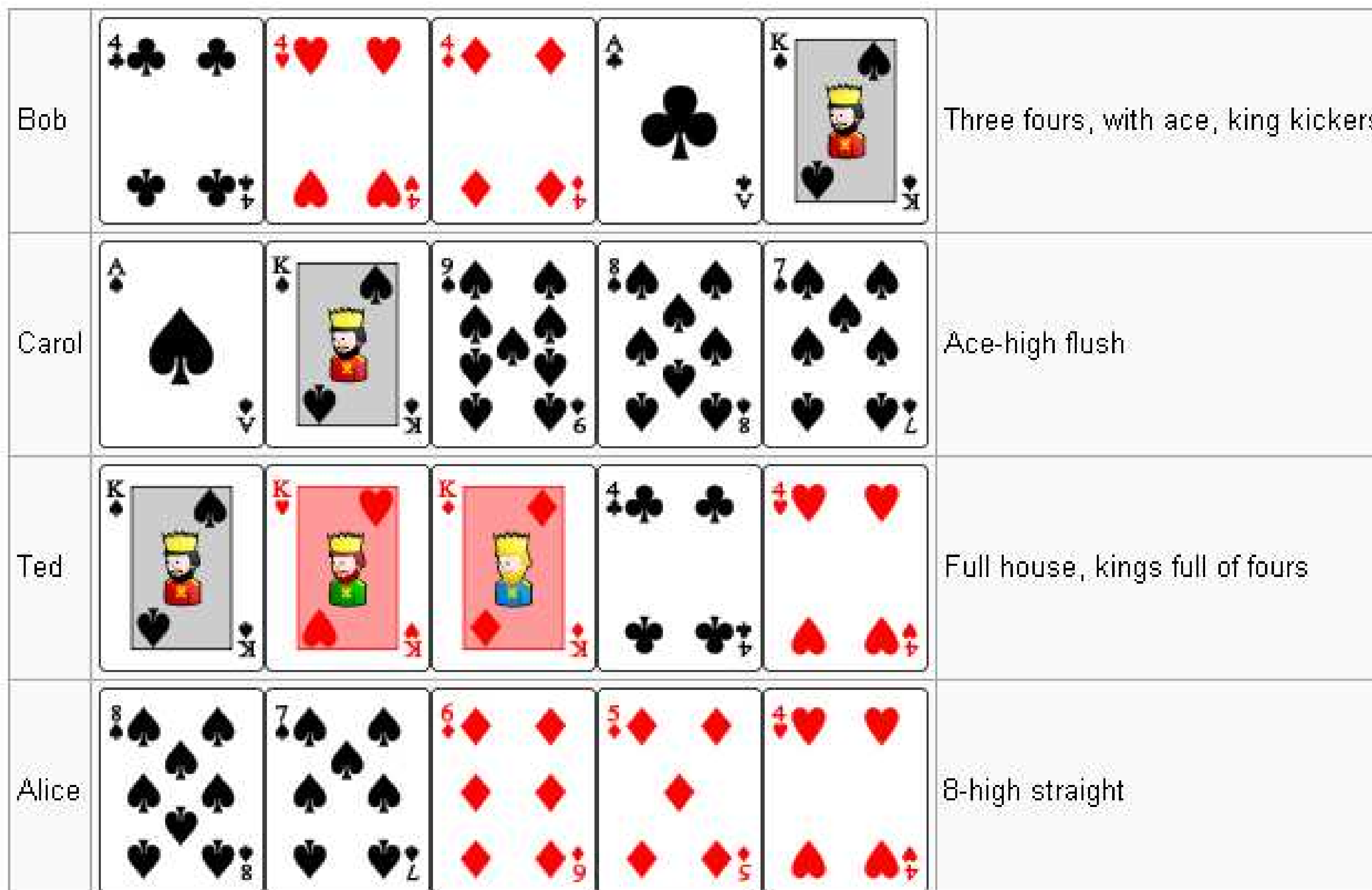




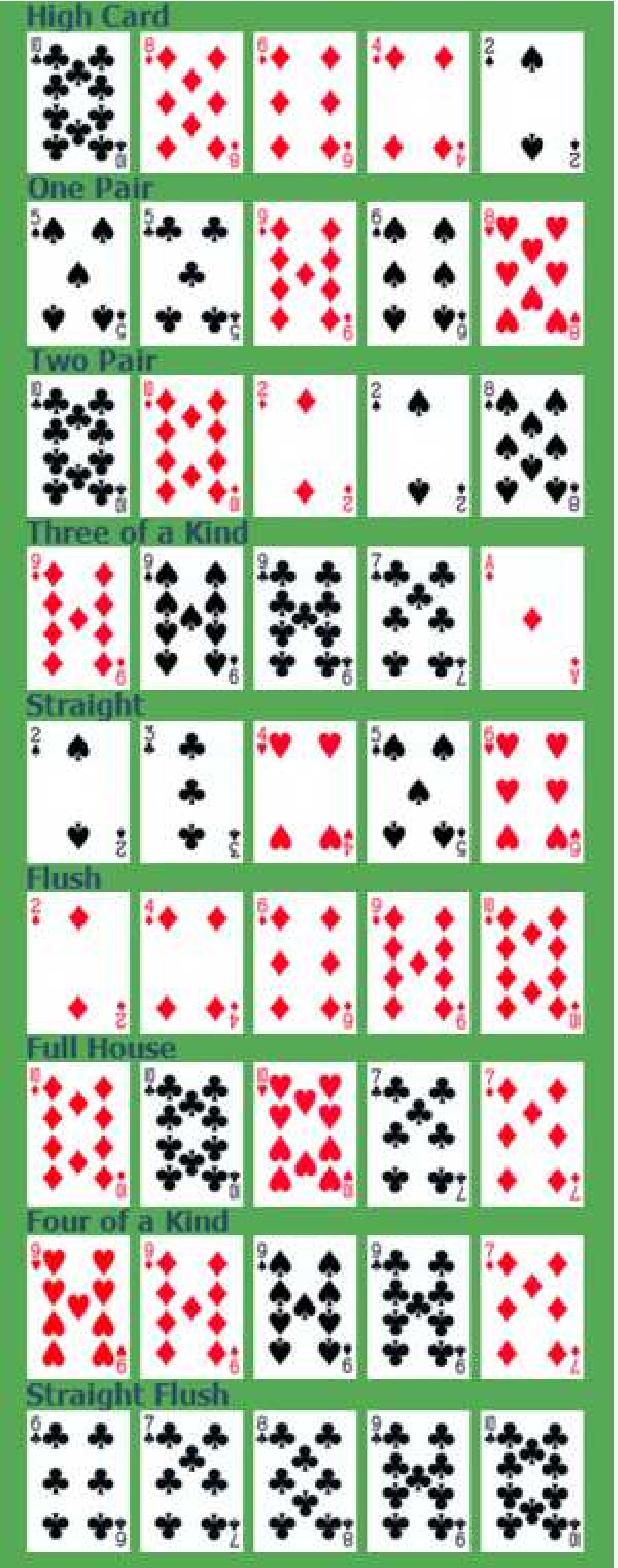
Sample showdown



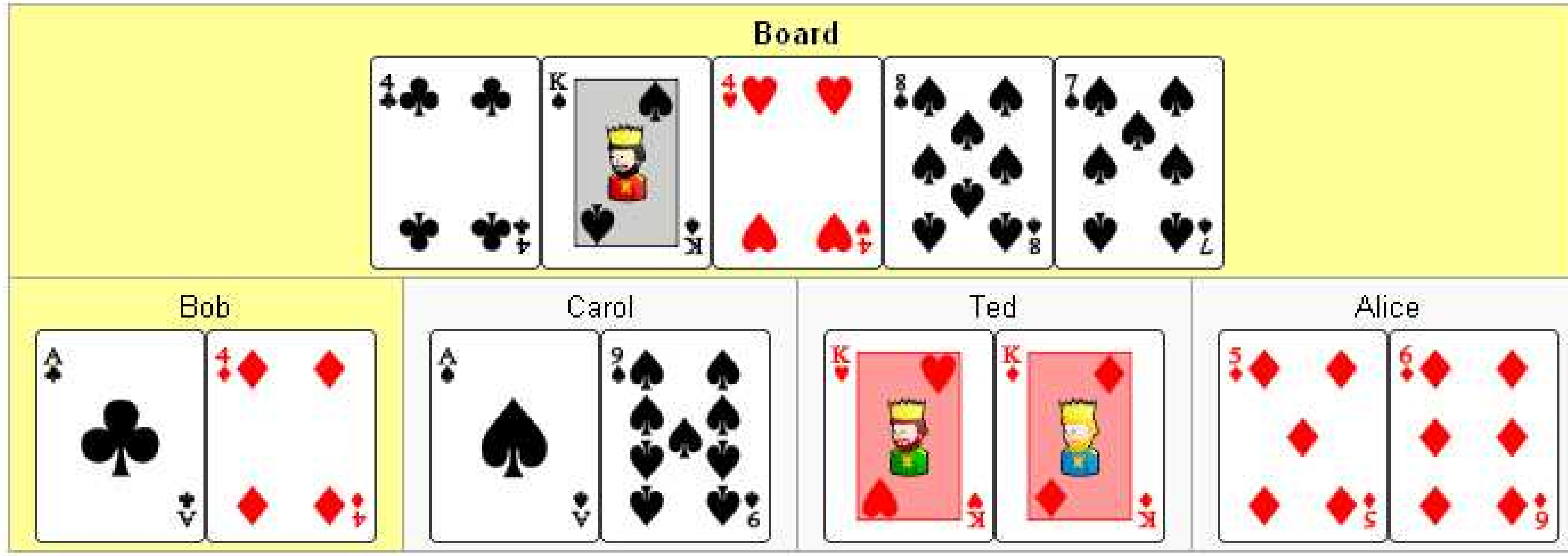
Each player plays the best 5-card hand they can make with the seven cards available. They have



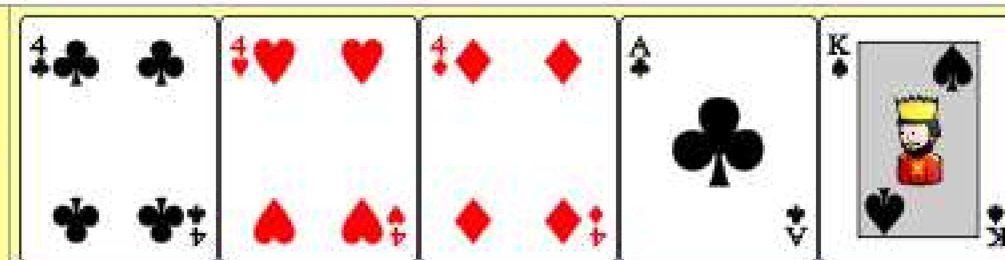
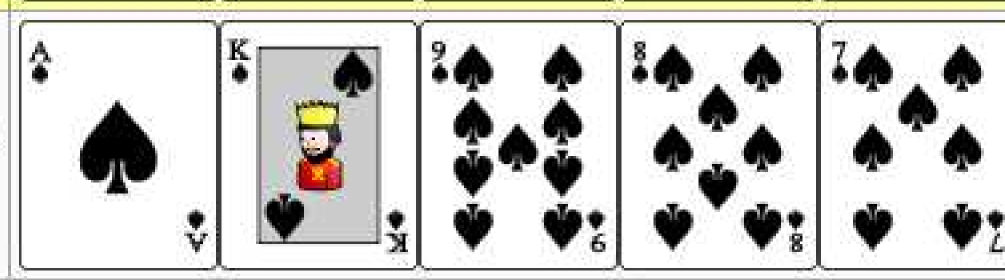
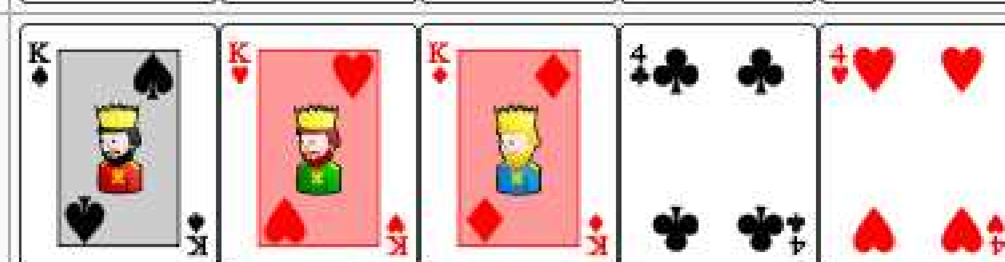
In this case, Ted's full house is the best hand, with Carol in 2nd, Alice in 3rd and Bob last.



Sample showdown



Each player plays the best 5-card hand they can make with the seven cards available. They have

Bob		Three fours, with ace, king kickers
Carol		Ace-high flush
Ted		Full house, kings full of fours
Alice		8-high straight

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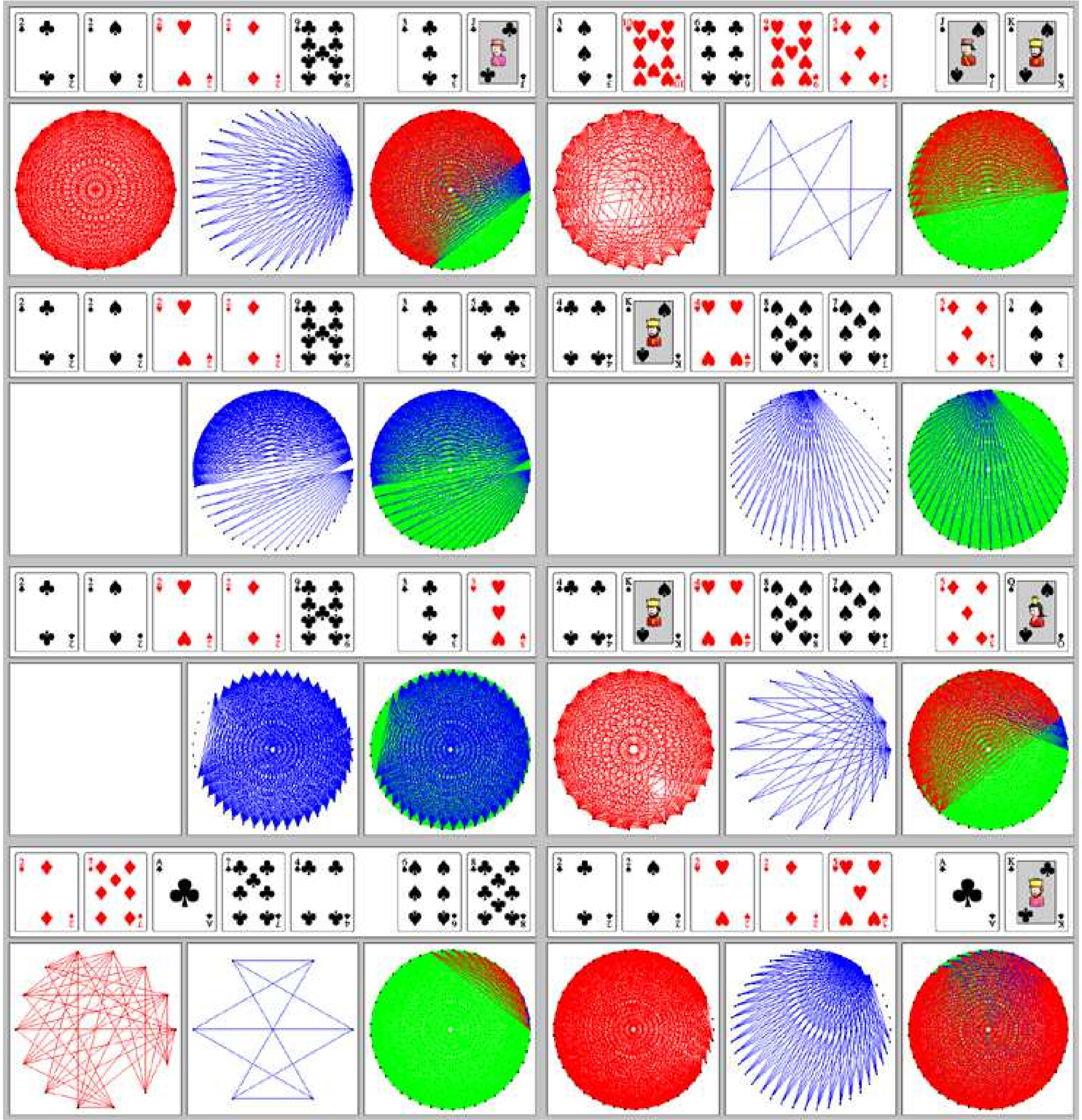

The algorithm

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Applicable to general simple graphs

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Calculates coefficients one at a time

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Calculates coefficients one at a time

Brute force algorithm

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Brute force algorithm

Extensive use of three main resources of computers

- Computational Power
- Main Memory (RAM)
- Hard Disc Space

The algorithm

Applicable to general simple graphs

Calculates coefficients one at a time

Brute force algorithm

Extensive use of three main resources of computers

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Achieve

- Maximal factorization within a single and among several graphs
- Reduction of computation time

Performance

Performance

"What experience I have suggests that it could be difficult to develop something useful for graphs on 12 vertices. So getting the first nine coefficients of just one graph on 45 vertices seems impossible at the moment." - Chris Godsil

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An algorithm for calculating the independence and vertex-cover polynomials of a graph
Applied Mathematics and Computation 190, 2007
Gordon G. Cash

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Computed first five coefficients of 151215 graphs on 45 vertices in less than one week using

- 4 x 2.67 GHz Quad-Core CPU
- 2 GB RAM
- 70 GB hard disc space

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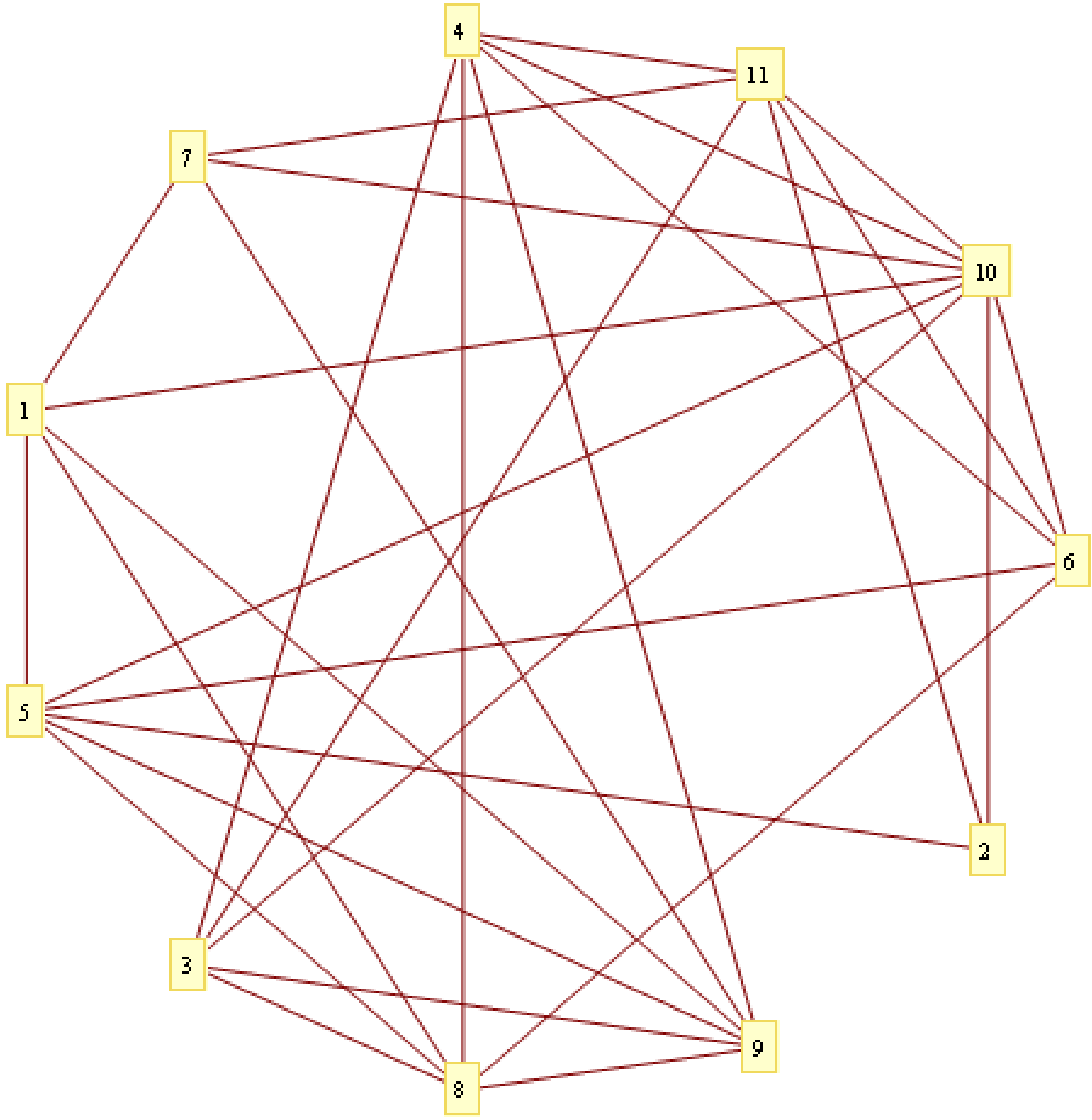
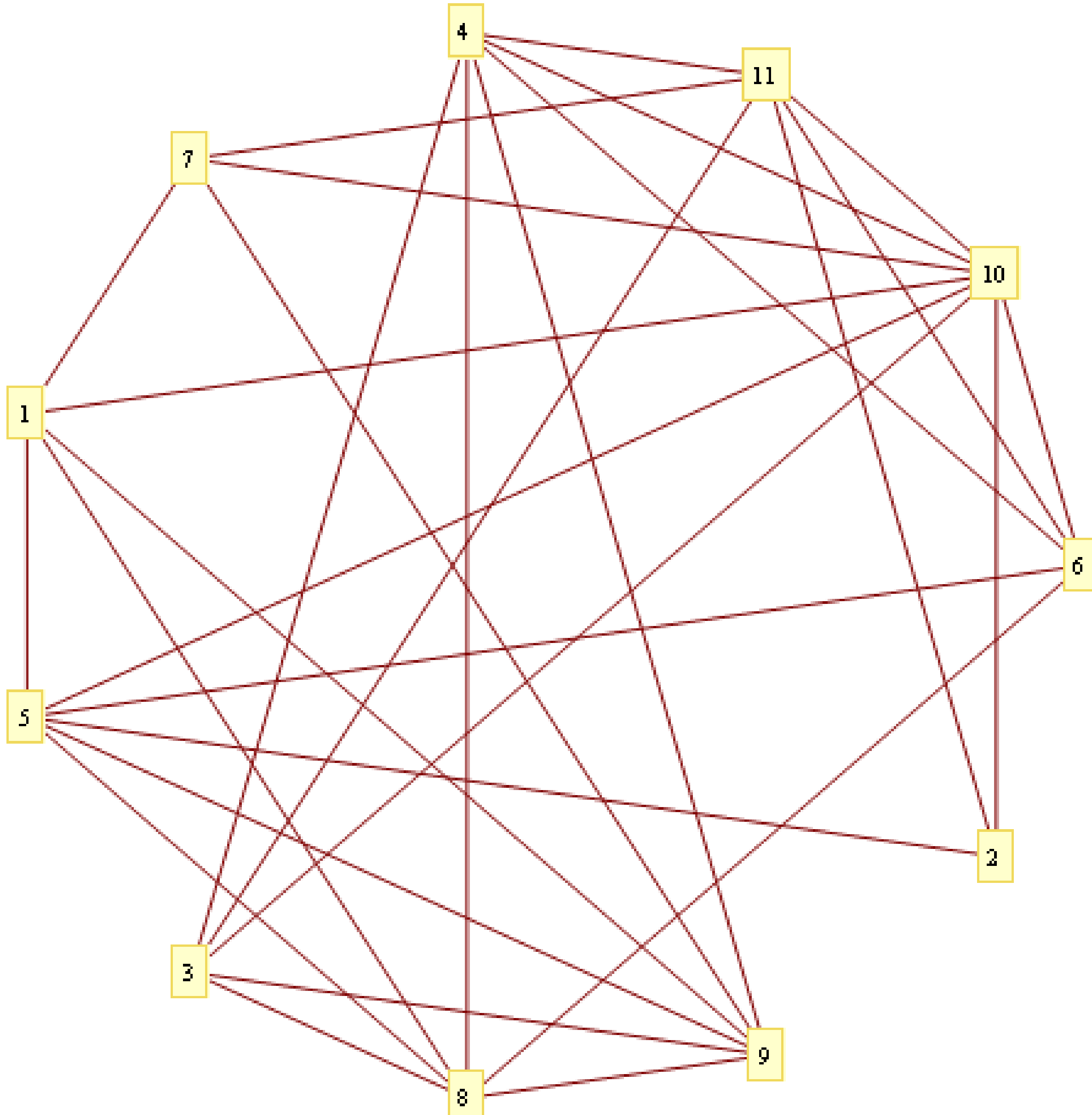
Open question: Complexity

Step 1: Preparation of graphs


4
11
7
10
1
6
5
2
3
9
8



6 → 10
10 → 11
10 → 4
7 → 1
5 → 1
3 → 4
6 → 4
7 → 10
10 → 5
5 → 8
8 → 9
7 → 9
10 → 3
11 → 7
2 → 10
1 → 8
2 → 5
5 → 6
11 → 6
2 → 11
4 → 11
10 → 1
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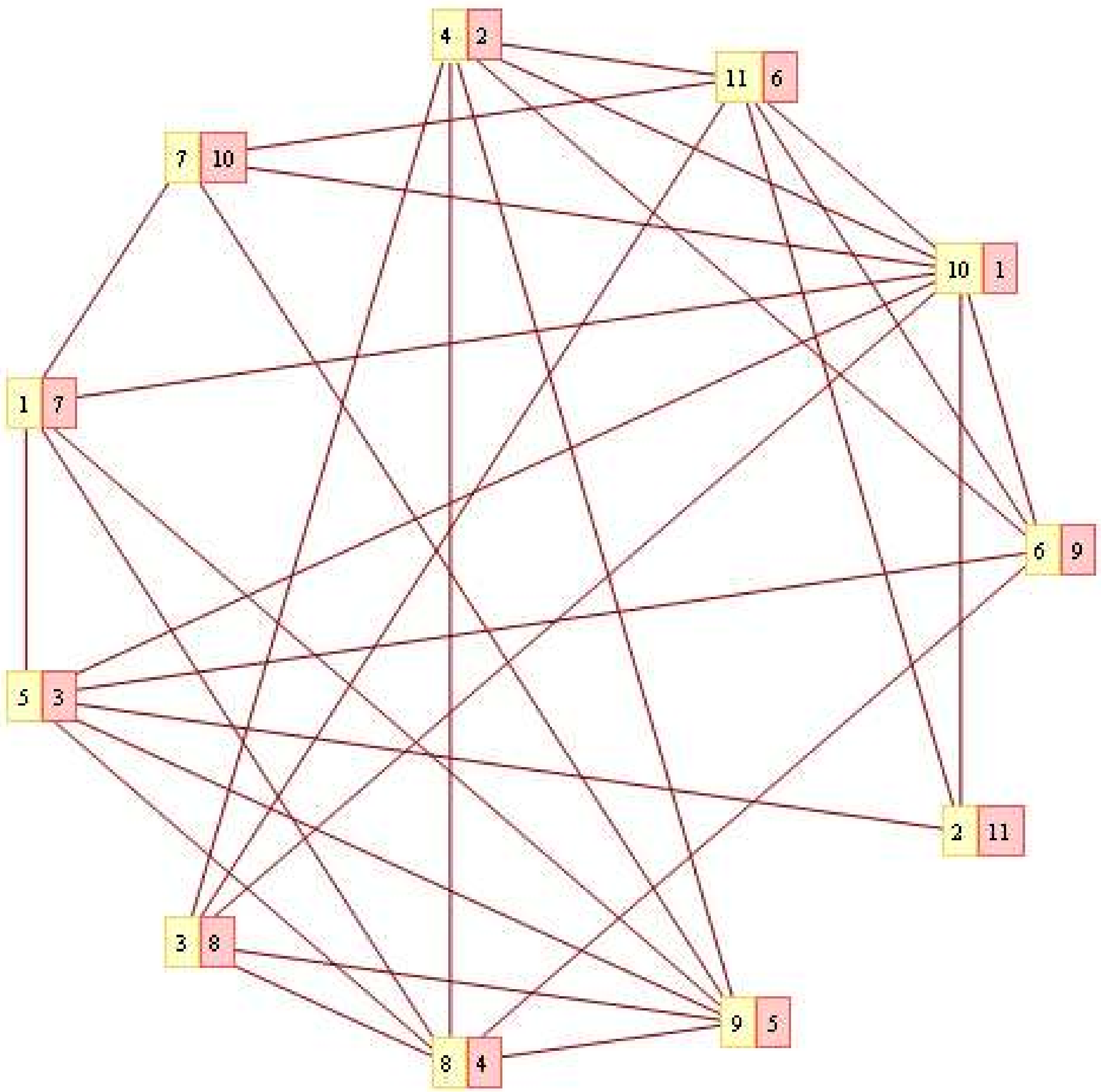
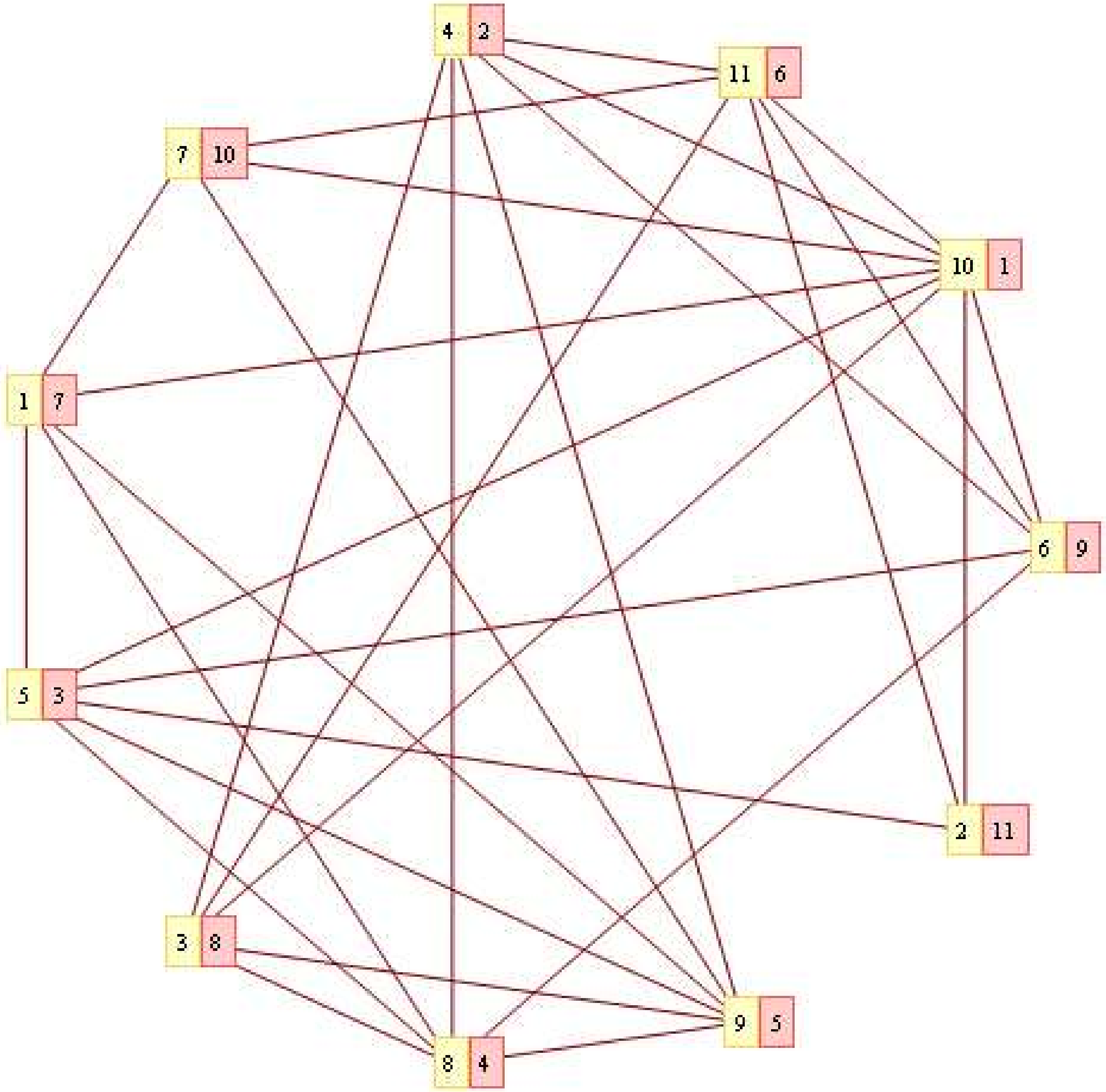
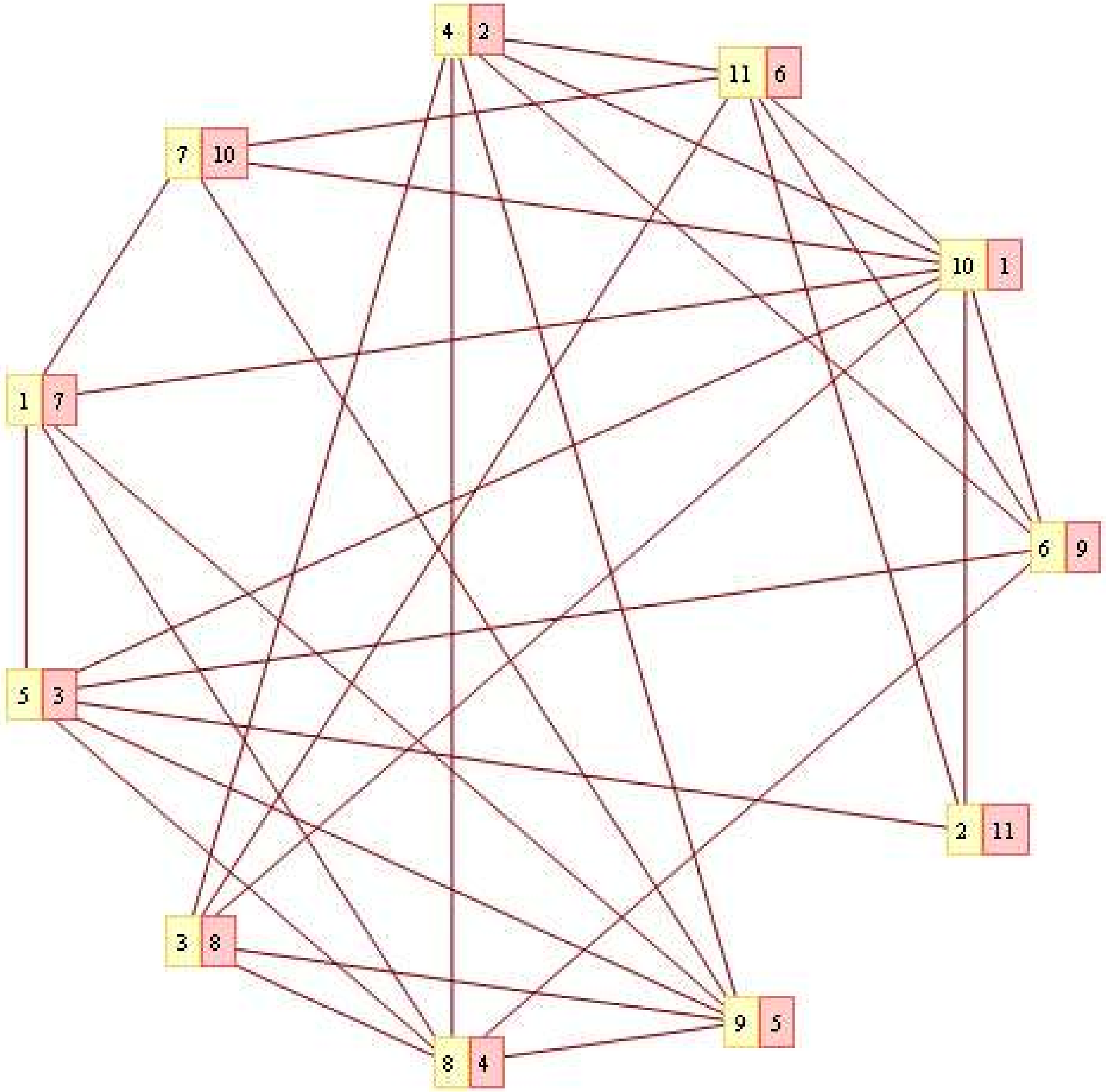
8 → 4
9 → 1
4 → 9
8 → 6
3 → 9
3 → 11
9 → 5



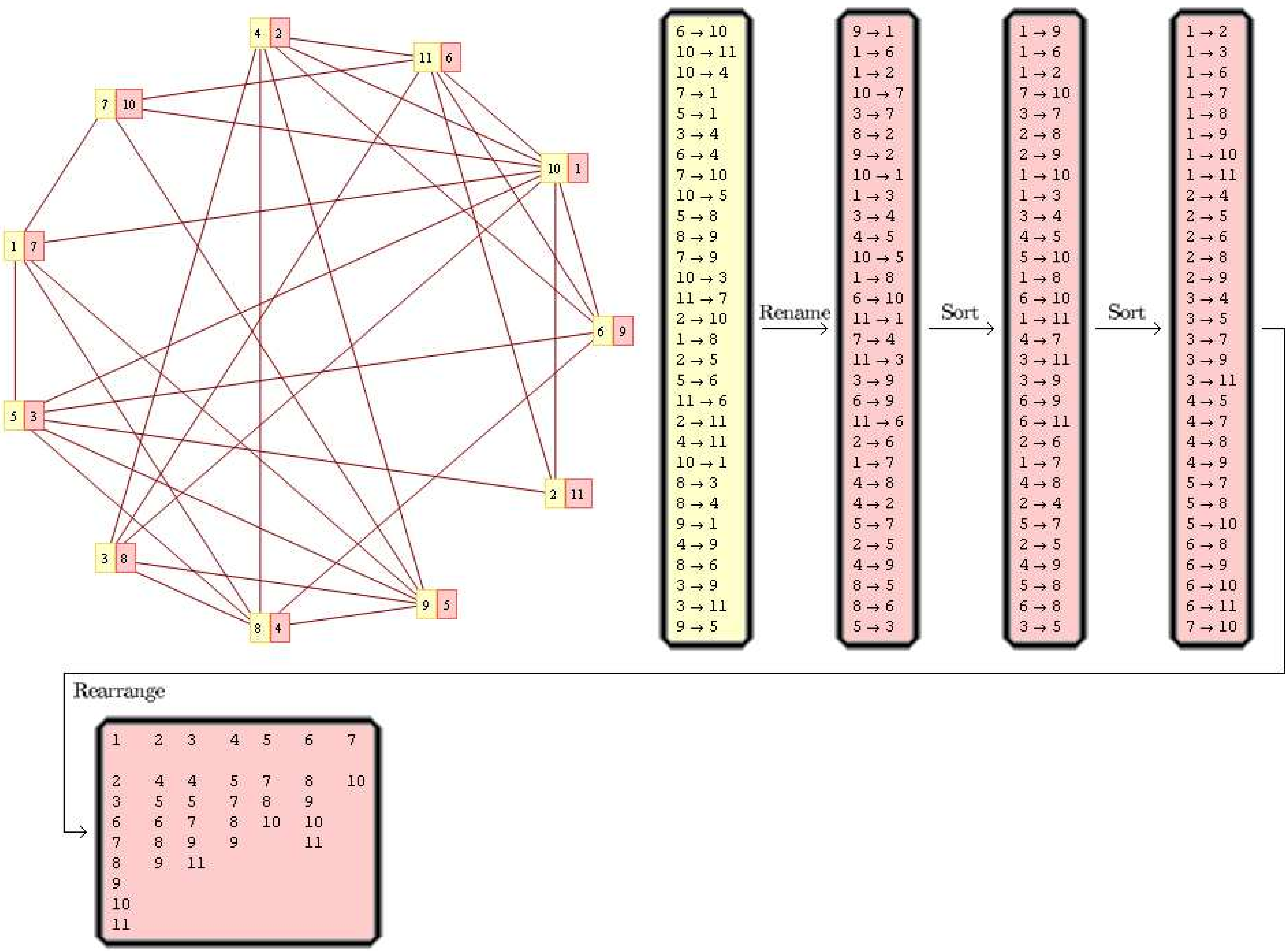
	Rename →	
6 → 10		9 → 1
10 → 11		1 → 6
10 → 4		1 → 2
7 → 1		10 → 7
5 → 1		3 → 7
3 → 4		8 → 2
6 → 4		9 → 2
7 → 10		10 → 1
10 → 5		1 → 3
5 → 8		3 → 4
8 → 9		4 → 5
7 → 9		10 → 5
10 → 3		1 → 8
11 → 7		6 → 10
2 → 10		11 → 1
1 → 8		7 → 4
2 → 5		11 → 3
5 → 6		3 → 9
11 → 6		6 → 9
2 → 11		11 → 6
4 → 11		2 → 6
10 → 1		1 → 7
8 → 3		4 → 8
8 → 4		4 → 2
9 → 1		5 → 7
4 → 9		2 → 5
8 → 6		4 → 9
3 → 9		8 → 5
3 → 11		8 → 6
9 → 5		5 → 3

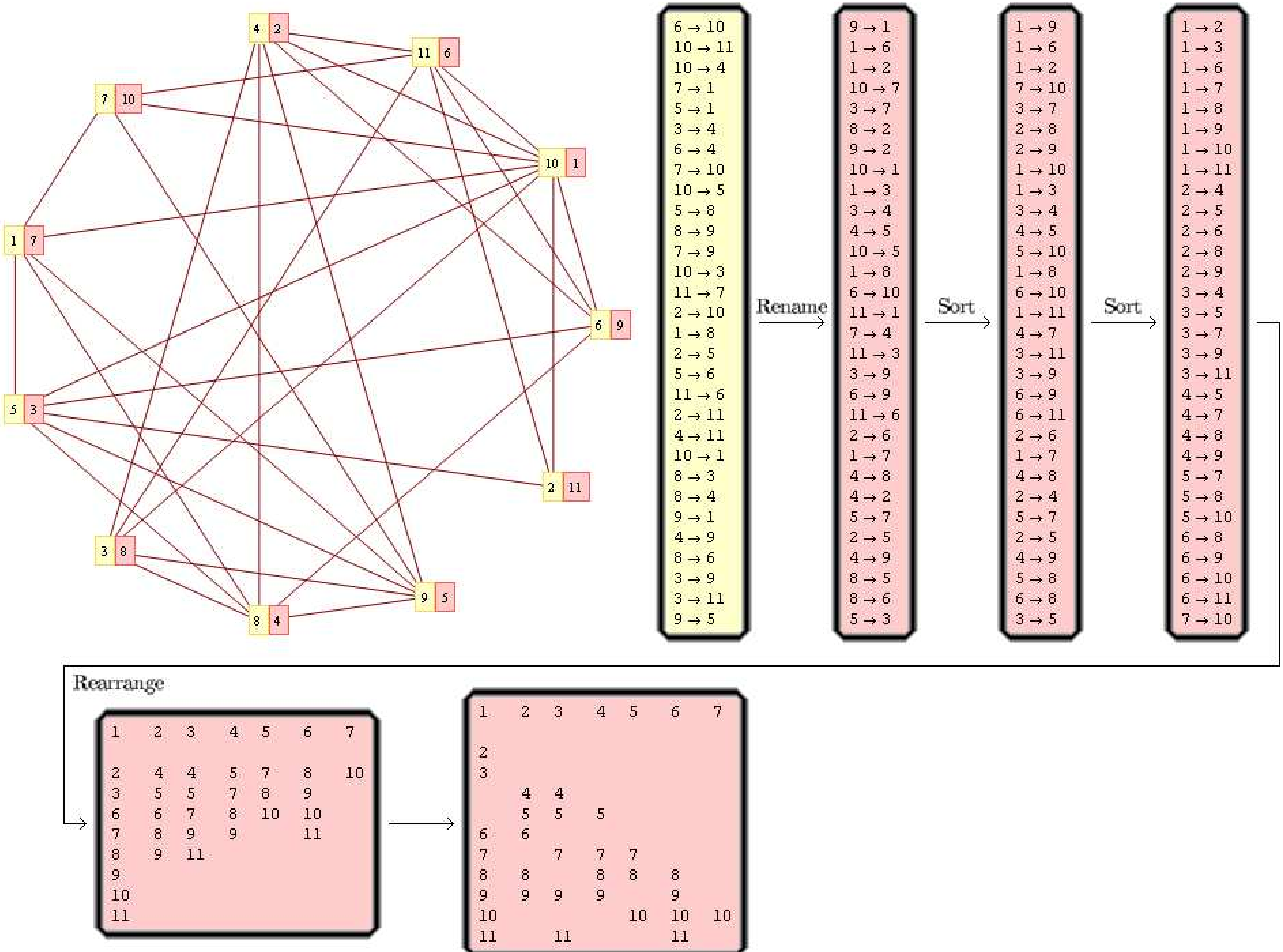


$6 \to 10$		$9 \to 1$		$1 \to 9$
$10 \to 11$		$1 \to 6$		$1 \to 6$
$10 \to 4$		$1 \to 2$		$1 \to 2$
$7 \to 1$		$10 \to 7$		$7 \to 10$
$5 \to 1$		$3 \to 7$		$3 \to 7$
$3 \to 4$		$8 \to 2$		$2 \to 8$
$6 \to 4$		$9 \to 2$		$2 \to 9$
$7 \to 10$		$10 \to 1$		$1 \to 10$
$10 \to 5$		$1 \to 3$		$1 \to 3$
$5 \to 8$		$3 \to 4$		$3 \to 4$
$8 \to 9$		$4 \to 5$		$4 \to 5$
$7 \to 9$		$10 \to 5$		$5 \to 10$
$10 \to 3$		$1 \to 8$		$1 \to 8$
$11 \to 7$		$6 \to 10$		$6 \to 10$
$2 \to 10$	Rename $\longrightarrow$	$11 \to 1$	Sort $\longrightarrow$	$1 \to 11$
$1 \to 8$		$7 \to 4$		$4 \to 7$
$2 \to 5$		$11 \to 3$		$3 \to 11$
$5 \to 6$		$3 \to 9$		$3 \to 9$
$11 \to 6$		$6 \to 9$		$6 \to 9$
$2 \to 11$		$11 \to 6$		$6 \to 11$
$4 \to 11$		$2 \to 6$		$2 \to 6$
$10 \to 1$		$1 \to 7$		$1 \to 7$
$8 \to 3$		$4 \to 8$		$4 \to 8$
$8 \to 4$		$4 \to 2$		$2 \to 4$
$9 \to 1$		$5 \to 7$		$5 \to 7$
$4 \to 9$		$2 \to 5$		$2 \to 5$
$8 \to 6$		$4 \to 9$		$4 \to 9$
$3 \to 9$		$8 \to 5$		$5 \to 8$
$3 \to 11$		$8 \to 6$		$6 \to 8$
$9 \to 5$		$5 \to 3$		$3 \to 5$



	Rename →	Sort →	Sort →
$6 \to 10$	$9 \to 1$	$1 \to 9$	$1 \to 2$
$10 \to 11$	$1 \to 6$	$1 \to 6$	$1 \to 3$
$10 \to 4$	$1 \to 2$	$1 \to 2$	$1 \to 6$
$7 \to 1$	$10 \to 7$	$7 \to 10$	$1 \to 7$
$5 \to 1$	$3 \to 7$	$3 \to 7$	$1 \to 8$
$3 \to 4$	$8 \to 2$	$2 \to 8$	$1 \to 9$
$6 \to 4$	$9 \to 2$	$2 \to 9$	$1 \to 10$
$7 \to 10$	$10 \to 1$	$1 \to 10$	$1 \to 11$
$10 \to 5$	$1 \to 3$	$1 \to 3$	$2 \to 4$
$5 \to 8$	$3 \to 4$	$3 \to 4$	$2 \to 5$
$8 \to 9$	$4 \to 5$	$4 \to 5$	$2 \to 6$
$7 \to 9$	$10 \to 5$	$5 \to 10$	$2 \to 8$
$10 \to 3$	$1 \to 8$	$1 \to 8$	$2 \to 9$
$11 \to 7$	$6 \to 10$	$6 \to 10$	$3 \to 4$
$2 \to 10$	$11 \to 1$	$1 \to 11$	$3 \to 5$
$1 \to 8$	$7 \to 4$	$4 \to 7$	$3 \to 7$
$2 \to 5$	$11 \to 3$	$3 \to 11$	$3 \to 9$
$5 \to 6$	$3 \to 9$	$3 \to 9$	$3 \to 11$
$11 \to 6$	$6 \to 9$	$6 \to 9$	$4 \to 5$
$2 \to 11$	$11 \to 6$	$6 \to 11$	$4 \to 7$
$4 \to 11$	$2 \to 6$	$2 \to 6$	$4 \to 8$
$10 \to 1$	$1 \to 7$	$1 \to 7$	$4 \to 9$
$8 \to 3$	$4 \to 8$	$4 \to 8$	$5 \to 7$
$8 \to 4$	$4 \to 2$	$2 \to 4$	$5 \to 8$
$9 \to 1$	$5 \to 7$	$5 \to 7$	$5 \to 10$
$4 \to 9$	$2 \to 5$	$2 \to 5$	$6 \to 8$
$8 \to 6$	$4 \to 9$	$4 \to 9$	$6 \to 9$
$3 \to 9$	$8 \to 5$	$5 \to 8$	$6 \to 10$
$3 \to 11$	$8 \to 6$	$6 \to 8$	$6 \to 11$
$9 \to 5$	$5 \to 3$	$3 \to 5$	$7 \to 10$


4 2
11 6
7 10
10 1
1 7
6 9
5 3
2 11
3 8
9 5
8 4
6 → 10
10 → 11
10 → 4
7 → 1
5 → 1
3 → 4
6 → 4
7 → 10
10 → 5
5 → 8
8 → 9
7 → 9
10 → 3
11 → 7
2 → 10
1 → 8
2 → 5
5 → 6
11 → 6
2 → 11
4 → 11
10 → 1
8 → 3
8 → 4
9 → 1
4 → 9
8 → 6
3 → 9
3 → 11
9 → 5
Rename
9 → 1
1 → 6
1 → 2
10 → 7
3 → 7
8 → 2
9 → 2
10 → 1
1 → 3
3 → 4
4 → 5
10 → 5
1 → 8
6 → 10
11 → 1
7 → 4
11 → 3
3 → 9
6 → 9
11 → 6
2 → 6
1 → 7
4 → 8
4 → 2
5 → 7
2 → 5
4 → 9
8 → 5
8 → 6
5 → 3
Sort
1 → 9
1 → 6
1 → 2
7 → 10
3 → 7
2 → 8
2 → 9
1 → 10
1 → 3
3 → 4
4 → 5
5 → 10
1 → 8
6 → 10
1 → 11
4 → 7
3 → 11
3 → 9
6 → 9
6 → 11
2 → 6
1 → 7
4 → 8
2 → 4
5 → 7
2 → 5
4 → 9
5 → 8
6 → 8
3 → 5
Sort
1 → 2
1 → 3
1 → 6
1 → 7
1 → 8
1 → 9
1 → 10
1 → 11
2 → 4
2 → 5
2 → 6
2 → 8
2 → 9
3 → 4
3 → 5
3 → 7
3 → 9
3 → 11
4 → 5
4 → 7
4 → 8
4 → 9
5 → 7
5 → 8
5 → 10
6 → 8
6 → 9
6 → 10
6 → 11
7 → 10
Rearrange
1 2 3 4 5 6 7
2 4 4 5 7 8 10
3 5 5 7 8 9
6 6 7 8 10 10
7 8 9 9 11
8 9 11
9
10
11


6 → 10
10 → 11
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7 → 1
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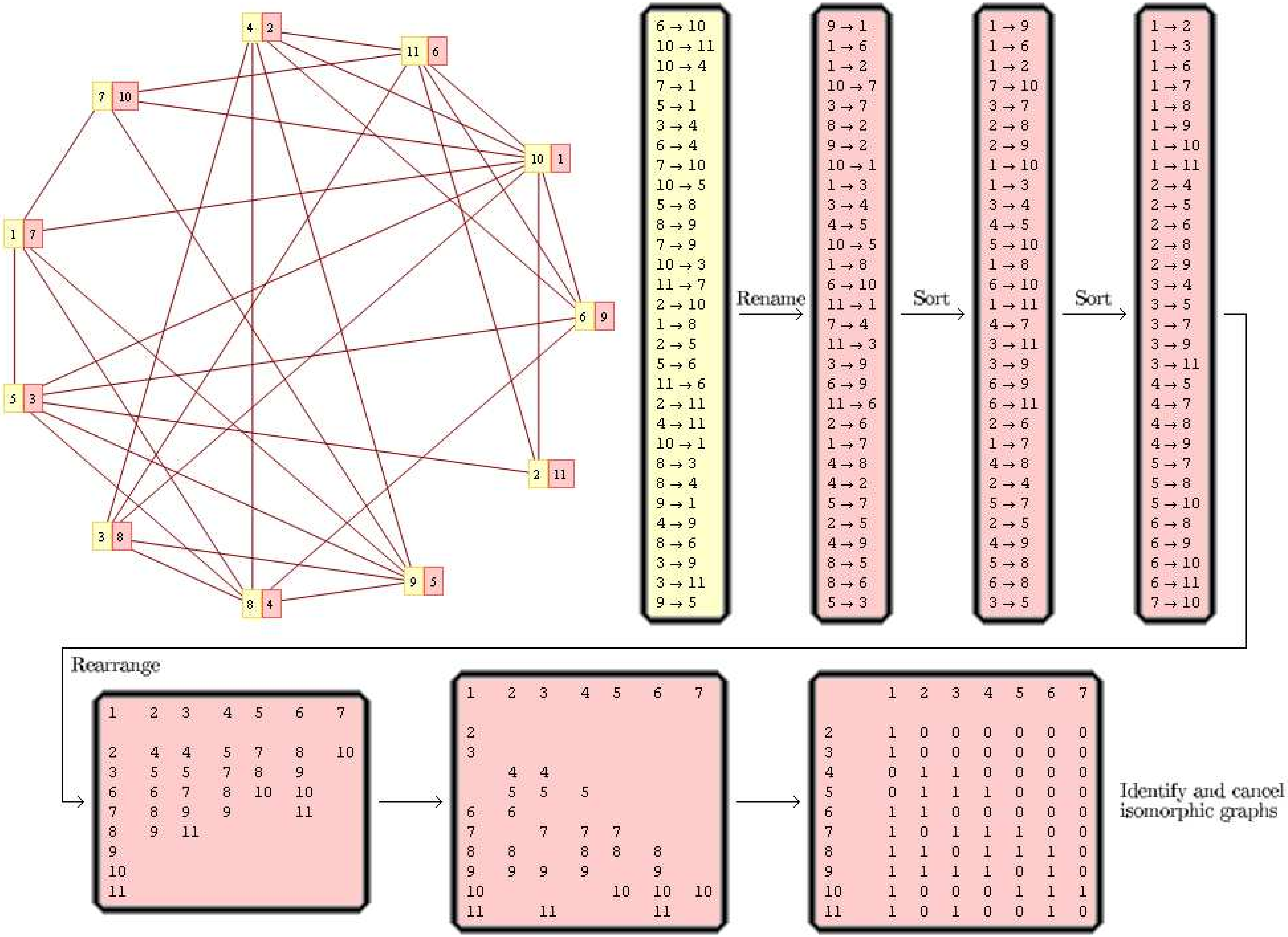
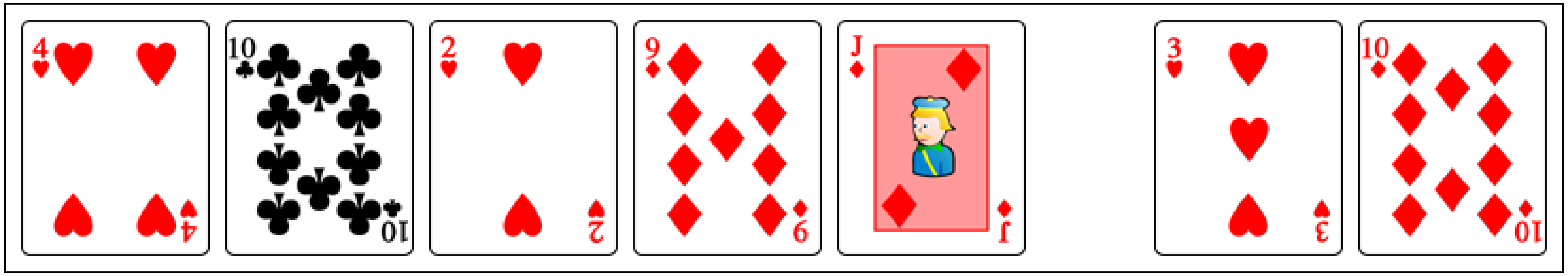
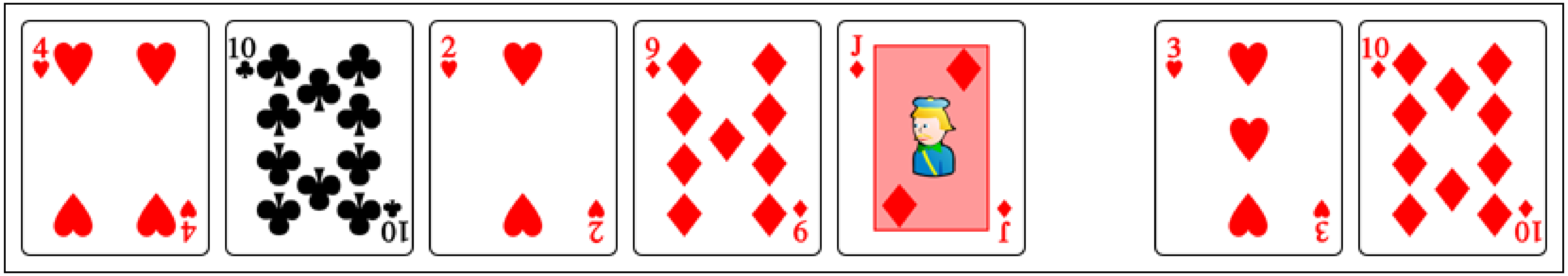
7 → 9
10 → 3
11 → 7
2 → 10
1 → 8
2 → 5
5 → 6
11 → 6
2 → 11
4 → 11
10 → 1
8 → 3
8 → 4
9 → 1
4 → 9
8 → 6
3 → 9
3 → 11
9 → 5
Rename
9 → 1
1 → 6
1 → 2
10 → 7
3 → 7
8 → 2
9 → 2
10 → 1
1 → 3
3 → 4
4 → 5
10 → 5
1 → 8
6 → 10
11 → 1
7 → 4
11 → 3
3 → 9
6 → 9
11 → 6
2 → 6
1 → 7
4 → 8
4 → 2
5 → 7
2 → 5
4 → 9
8 → 5
8 → 6
5 → 3
Sort
1 → 9
1 → 6
1 → 2
7 → 10
3 → 7
2 → 8
2 → 9
1 → 10
1 → 3
3 → 4
4 → 5
5 → 10
1 → 8
6 → 10
1 → 11
4 → 7
3 → 11
3 → 9
6 → 9
6 → 11
2 → 6
1 → 7
4 → 8
2 → 4
5 → 7
2 → 5
4 → 9
5 → 8
6 → 8
3 → 5
Sort
1 → 2
1 → 3
1 → 6
1 → 7
1 → 8
1 → 9
1 → 10
1 → 11
2 → 4
2 → 5
2 → 6
2 → 8
2 → 9
3 → 4
3 → 5
3 → 7
3 → 9
3 → 11
4 → 5
4 → 7
4 → 8
4 → 9
5 → 7
5 → 8
5 → 10
6 → 8
6 → 9
6 → 10
6 → 11
7 → 10
Rearrange
1 2 3 4 5 6 7
2 4 4 5 7 8 10
3 5 5 7 8 9
6 6 7 8 10 10
7 8 9 9 11
8 9 11
9
10
11
1 2 3 4 5 6 7
2
3
4 4
5 5 5
6 6
7 7 7 7
8 8 8 8 8
9 9 9 9 9
10 10 10 10
11 11 11

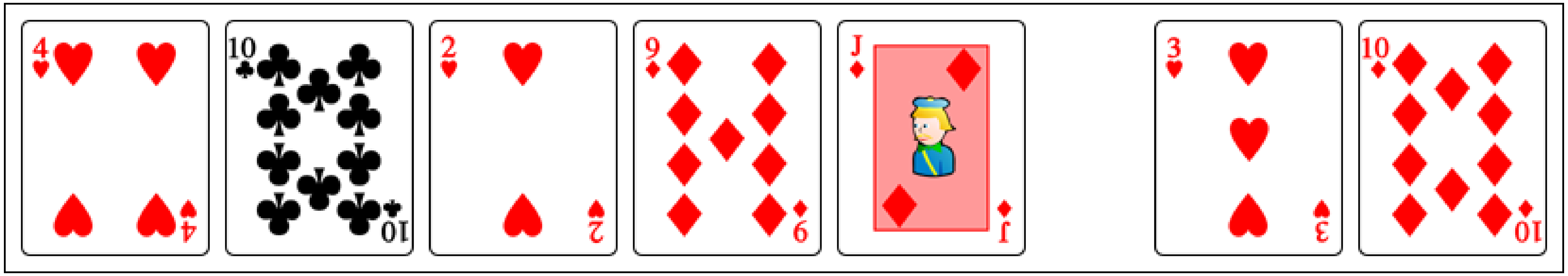

4 2
11 6
7 10
10 1
1 7
6 9
5 3
2 11
3 8
9 5
8 4
6 → 10, 10 → 11, 10 → 4, 7 → 1, 5 → 1, 3 → 4, 6 → 4, 7 → 10, 10 → 5, 5 → 8, 8 → 9, 7 → 9, 10 → 3, 11 → 7, 2 → 10, 1 → 8, 2 → 5, 5 → 6, 11 → 6, 2 → 11, 4 → 11, 10 → 1, 8 → 3, 8 → 4, 9 → 1, 4 → 9, 8 → 6, 3 → 9, 3 → 11, 9 → 5
Rename
9 → 1, 1 → 6, 1 → 2, 10 → 7, 3 → 7, 8 → 2, 9 → 2, 10 → 1, 1 → 3, 3 → 4, 4 → 5, 10 → 5, 1 → 8, 6 → 10, 11 → 1, 7 → 4, 11 → 3, 3 → 9, 6 → 9, 11 → 6, 2 → 6, 1 → 7, 4 → 8, 4 → 2, 5 → 7, 2 → 5, 4 → 9, 8 → 5, 8 → 6, 5 → 3
Sort
1 → 9, 1 → 6, 1 → 2, 7 → 10, 3 → 7, 2 → 8, 2 → 9, 1 → 10, 1 → 3, 3 → 4, 4 → 5, 5 → 10, 1 → 8, 6 → 10, 1 → 11, 4 → 7, 3 → 11, 3 → 9, 6 → 9, 6 → 11, 2 → 6, 1 → 7, 4 → 8, 2 → 4, 5 → 7, 2 → 5, 4 → 9, 5 → 8, 6 → 8, 3 → 5
Sort
1 → 2, 1 → 3, 1 → 6, 1 → 7, 1 → 8, 1 → 9, 1 → 10, 1 → 11, 2 → 4, 2 → 5, 2 → 6, 2 → 8, 2 → 9, 3 → 4, 3 → 5, 3 → 7, 3 → 9, 3 → 11, 4 → 5, 4 → 7, 4 → 8, 4 → 9, 5 → 7, 5 → 8, 5 → 10, 6 → 8, 6 → 9, 6 → 10, 6 → 11, 7 → 10
Rearrange
1 2 3 4 5 6 7
2 4 4 5 7 8 10
3 5 5 7 8 9
6 6 7 8 10 10
7 8 9 9 11
8 9 11
9
10
11
1 2 3 4 5 6 7
2
3
4 4
5 5 5
6 6
7 7 7 7
8 8 8 8 8
9 9 9 9 9
10 10 10 10
11 11 11
1 2 3 4 5 6 7
2 1 0 0 0 0 0 0
3 1 0 0 0 0 0 0
4 0 1 1 0 0 0 0
5 0 1 1 1 0 0 0
6 1 1 0 0 0 0 0
7 1 0 1 1 1 0 0
8 1 1 0 1 1 1 0
9 1 1 1 1 0 1 0
10 1 0 0 0 1 1 1
11 1 0 1 0 0 1 0
Identify and cancel isomorphic graphs

How many graphs occur in Texas Hold'em Poker?



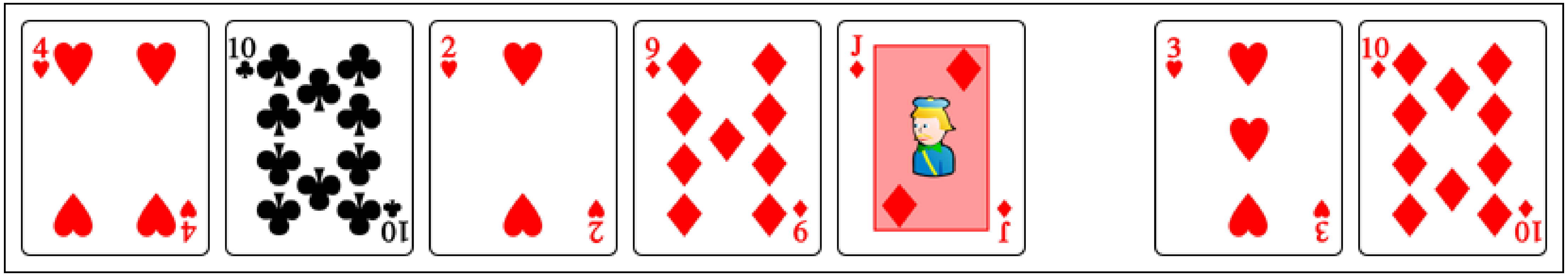


$$52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 \cdot 47 \cdot 46 = \underline{674274182400}$$



$$52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 \cdot 47 \cdot 46 = \underline{674274182400}$$

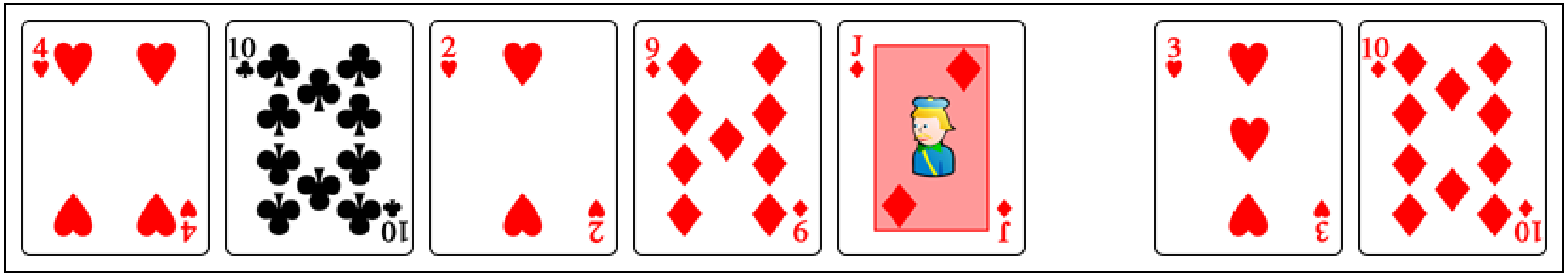
$$\binom{52}{5} \cdot \binom{47}{2} = \underline{2809475760}$$



$$52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 \cdot 47 \cdot 46 = \underline{674274182400}$$

$$\binom{52}{5} \cdot \binom{47}{2} = \underline{2809475760}$$

$$558883 \cdot 715 = \underline{399601345}$$

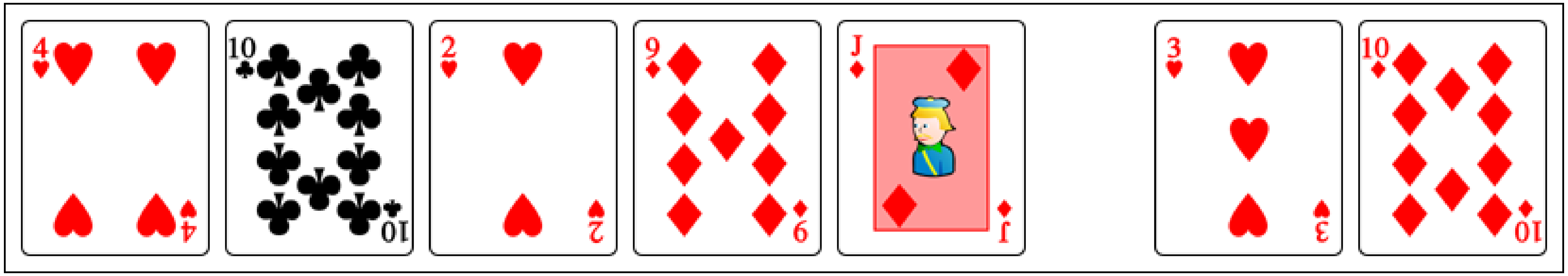


$$52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 \cdot 47 \cdot 46 = \underline{674274182400}$$

$$\binom{52}{5} \cdot \binom{47}{2} = \underline{2809475760}$$

$$558883 \cdot 715 = \underline{399601345}$$

$$\underline{157506622}$$



$$52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 \cdot 47 \cdot 46 = \underline{674274182400}$$

$$\binom{52}{5} \cdot \binom{47}{2} = \underline{2809475760}$$

$$558883 \cdot 715 = \underline{399601345}$$

$$\underline{157506622}$$

$$\underline{\underline{151215}}$$

Step 2: Factorization

	1	2	3	4	5	6	7
2	1	0	0	0	0	0	0
3	1	0	0	0	0	0	0
4	0	1	1	0	0	0	0
5	0	1	1	1	0	0	0
6	1	1	0	0	0	0	0
7	1	0	1	1	1	0	0
8	1	1	0	1	1	1	0
9	1	1	1	1	0	1	0
10	1	0	0	0	1	1	1
11	1	0	1	0	0	1	0

	1	2	3	4	5	6	7
2	1	0	0	0	0	0	0
3	1	0	0	0	0	0	0
4	0	1	1	0	0	0	0
5	0	1	1	1	0	0	0
6	1	1	0	0	0	0	0
7	1	0	1	1	1	0	0
8	1	1	0	1	1	1	0
9	1	1	1	1	0	1	0
10	1	0	0	0	1	1	1
11	1	0	1	0	0	1	0

Select

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
4	0	0	0	0
5	0	1	0	0
6	1	0	0	0
7	1	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

	1	2	3	4	5	6	7
2	1	0	0	0	0	0	0
3	1	0	0	0	0	0	0
4	0	1	1	0	0	0	0
5	0	1	1	1	0	0	0
6	1	1	0	0	0	0	0
7	1	0	1	1	1	0	0
8	1	1	0	1	1	1	0
9	1	1	1	1	0	1	0
10	1	0	0	0	1	1	1
11	1	0	1	0	0	1	0

Select

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
4	0	0	0	0
5	0	1	0	0
6	1	0	0	0
7	1	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

↓Cancel

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
5	0	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

	1	2	3	4	5	6	7
2	1	0	0	0	0	0	0
3	1	0	0	0	0	0	0
4	0	1	1	0	0	0	0
5	0	1	1	1	0	0	0
6	1	1	0	0	0	0	0
7	1	0	1	1	1	0	0
8	1	1	0	1	1	1	0
9	1	1	1	1	0	1	0
10	1	0	0	0	1	1	1
11	1	0	1	0	0	1	0

Select

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
4	0	0	0	0
5	0	1	0	0
6	1	0	0	0
7	1	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

↓Cancel

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
5	0	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

1	0	0	0
1	0	0	0
0	1	0	0
1	1	1	0
1	1	1	0
1	0	1	1
1	0	1	0

	1	2	3	4	5	6	7
2	1	0	0	0	0	0	0
3	1	0	0	0	0	0	0
4	0	1	1	0	0	0	0
5	0	1	1	1	0	0	0
6	1	1	0	0	0	0	0
7	1	0	1	1	1	0	0
8	1	1	0	1	1	1	0
9	1	1	1	1	0	1	0
10	1	0	0	0	1	1	1
11	1	0	1	0	0	1	0

Select

	1	4	6	7
2	1	0	0	0
3	1	0	0	0
4	0	0	0	0
5	0	1	0	0
6	1	0	0	0
7	1	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

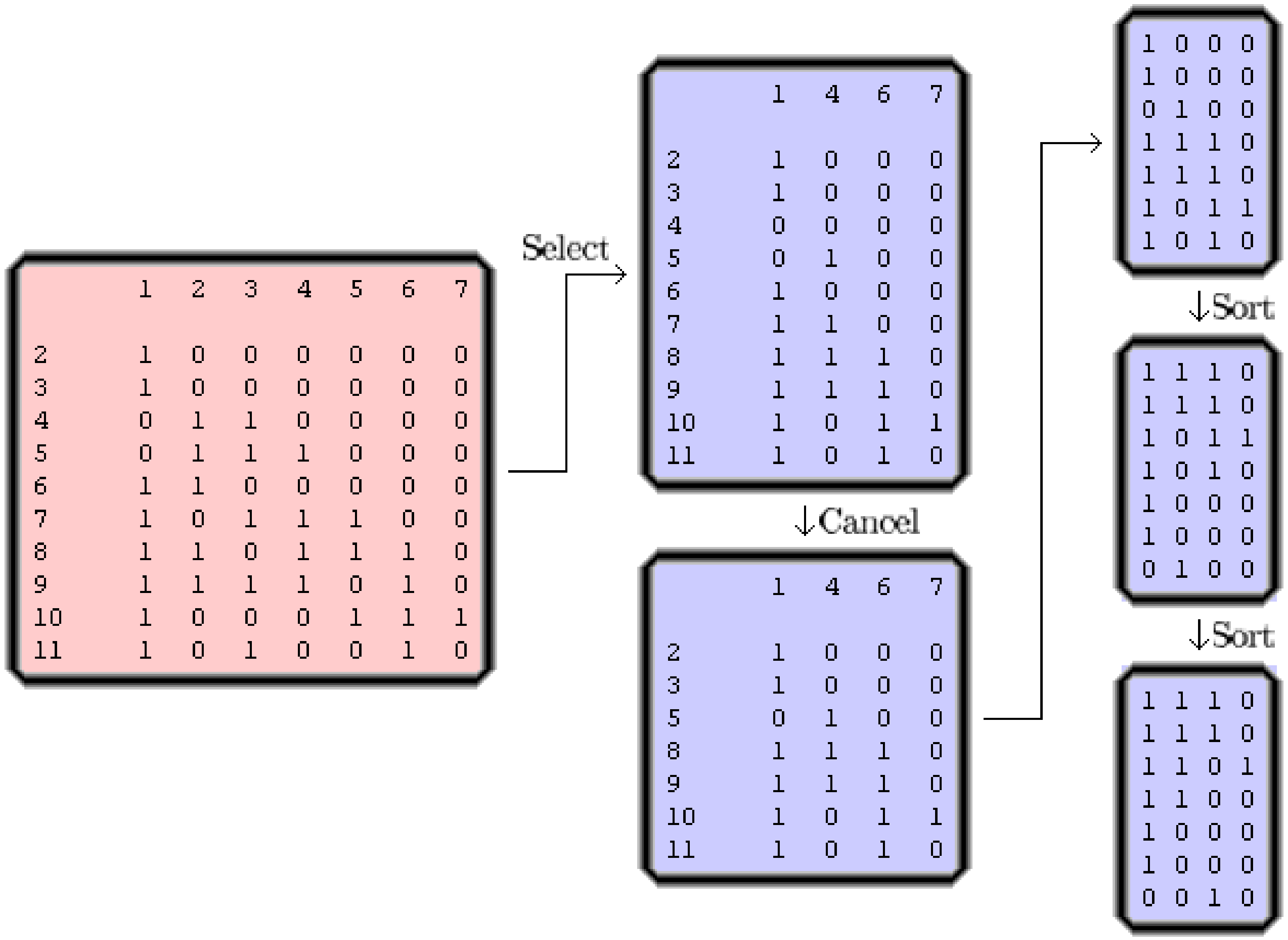
↓Cancel

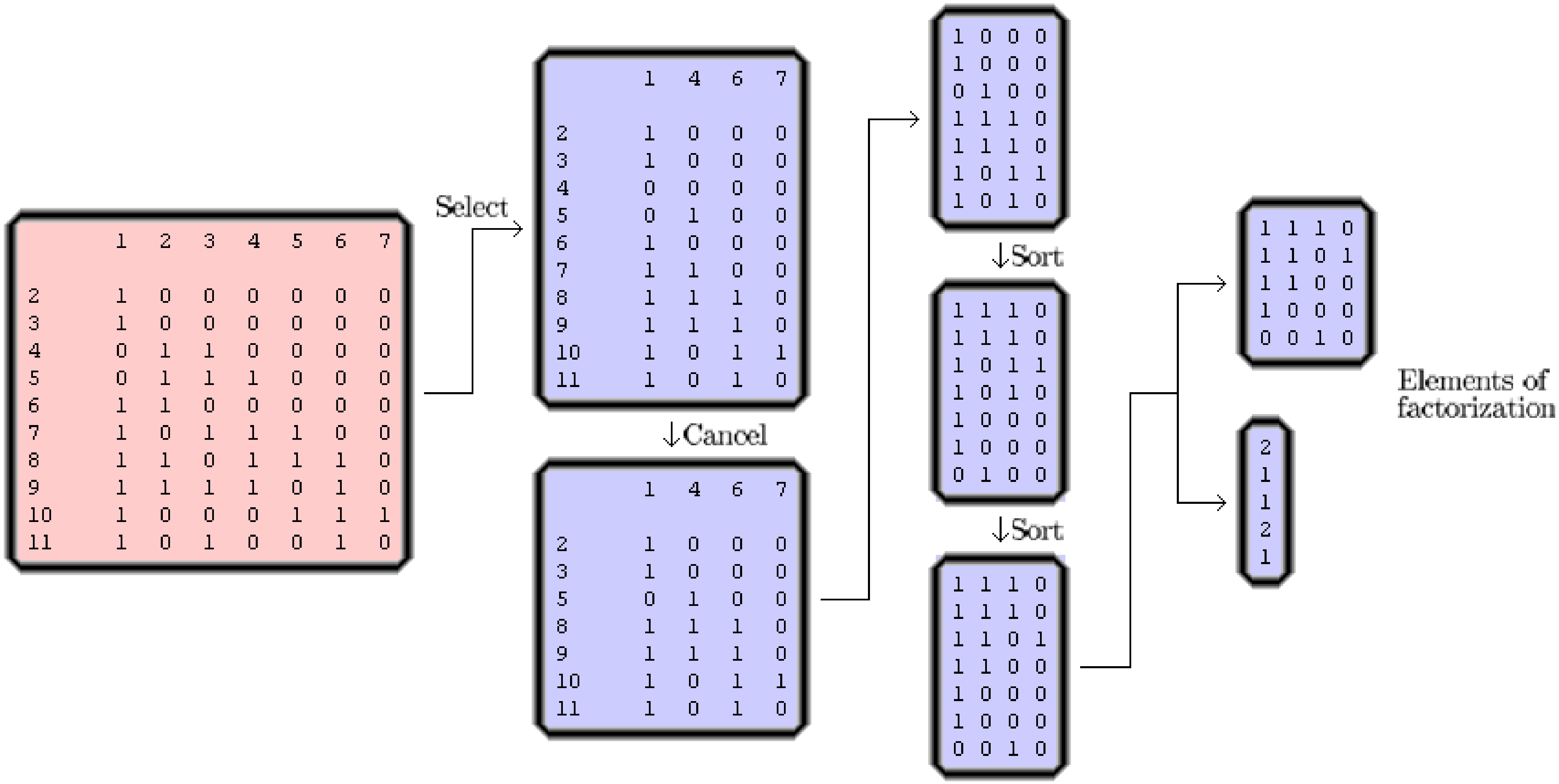
	1	4	6	7
2	1	0	0	0
3	1	0	0	0
5	0	1	0	0
8	1	1	1	0
9	1	1	1	0
10	1	0	1	1
11	1	0	1	0

1	0	0	0
1	0	0	0
0	1	0	0
1	1	1	0
1	1	1	0
1	0	1	1
1	0	1	0

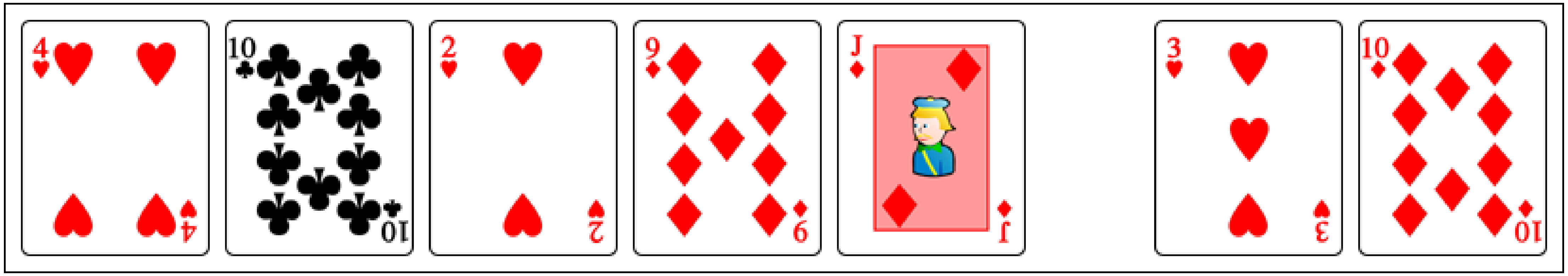
↓Sort

1	1	1	0
1	1	1	0
1	0	1	1
1	0	1	0
1	0	0	0
1	0	0	0
0	1	0	0


1 2 3 4 5 6 7
2 1 0 0 0 0 0 0
3 1 0 0 0 0 0 0
4 0 1 1 0 0 0 0
5 0 1 1 1 0 0 0
6 1 1 0 0 0 0 0
7 1 0 1 1 1 0 0
8 1 1 0 1 1 1 0
9 1 1 1 1 0 1 0
10 1 0 0 0 1 1 1
11 1 0 1 0 0 1 0
Select
1 4 6 7
2 1 0 0 0
3 1 0 0 0
4 0 0 0 0
5 0 1 0 0
6 1 0 0 0
7 1 1 0 0
8 1 1 1 0
9 1 1 1 0
10 1 0 1 1
11 1 0 1 0
↓Cancel
1 4 6 7
2 1 0 0 0
3 1 0 0 0
5 0 1 0 0
8 1 1 1 0
9 1 1 1 0
10 1 0 1 1
11 1 0 1 0
1 0 0 0
1 0 0 0
0 1 0 0
1 1 1 0
1 1 1 0
1 0 1 1
1 0 1 0
↓Sort
1 1 1 0
1 1 1 0
1 0 1 1
1 0 1 0
1 0 0 0
1 0 0 0
0 1 0 0
↓Sort
1 1 1 0
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
1 0 0 0
0 0 1 0

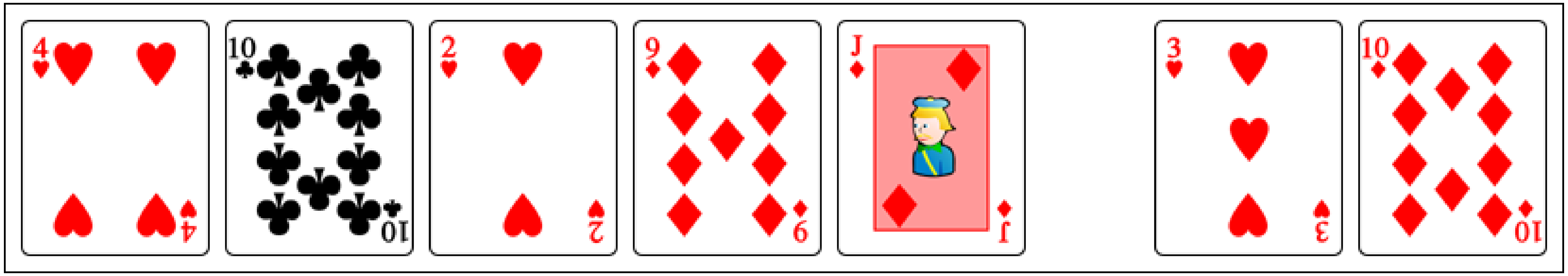


1 2 3 4 5 6 7
2 1 0 0 0 0 0 0
3 1 0 0 0 0 0 0
4 0 1 1 0 0 0 0
5 0 1 1 1 0 0 0
6 1 1 0 0 0 0 0
7 1 0 1 1 1 0 0
8 1 1 0 1 1 1 0
9 1 1 1 1 0 1 0
10 1 0 0 0 1 1 1
11 1 0 1 0 0 1 0
Select
1 4 6 7
2 1 0 0 0
3 1 0 0 0
4 0 0 0 0
5 0 1 0 0
6 1 0 0 0
7 1 1 0 0
8 1 1 1 0
9 1 1 1 0
10 1 0 1 1
11 1 0 1 0
↓Cancel
1 4 6 7
2 1 0 0 0
3 1 0 0 0
5 0 1 0 0
8 1 1 1 0
9 1 1 1 0
10 1 0 1 1
11 1 0 1 0
1 0 0 0
1 0 0 0
0 1 0 0
1 1 1 0
1 1 1 0
1 0 1 1
1 0 1 0
↓Sort
1 1 1 0
1 1 1 0
1 0 1 1
1 0 1 0
1 0 0 0
1 0 0 0
0 1 0 0
↓Sort
1 1 1 0
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
1 0 0 0
0 0 1 0
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
0 0 1 0
2
1
1
2
1
Elements of factorization



151215 Graphs

62593054397 Matrices

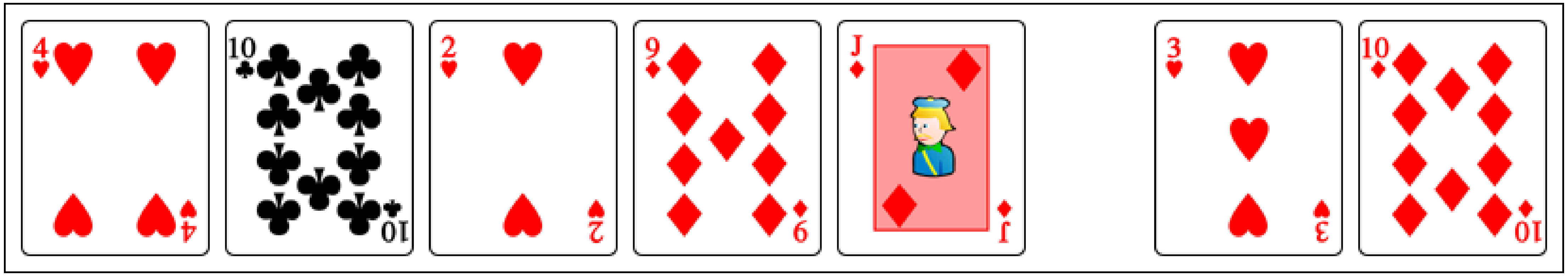


151215 Graphs

62593054397 Matrices

Reduce to:

333619695 Matrices



151215 Graphs

62593054397 Matrices

Reduce to:

333619695 Matrices

A factor of: 188

Step 3: Evaluation of matrices

$$\sum_{(i_1,\ldots,i_n)\in\{1,\ldots,m'\}^n} \left(\prod_{j=1}^{n} a'_{i_j,j} \left(b_{i_j} - \#\left\{ j' \in \{1,\ldots,j-1\} \,|\, i_{j'} = i_j \right\} \right) \right)$$

Ways to select one "1" in every column
of A such that no two "1"'s share a row

$$A = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$A' = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$b = \begin{pmatrix} 2 \\ 1 \\ 1 \\ 2 \\ 1 \end{pmatrix}$$

$$\sum_{(i_1,\ldots,i_n)\in\{1,\ldots,m'\}^n} \left(\prod_{j=1}^{n} a'_{i_j,j} \left(b_{i_j} - \#\left\{ j' \in \{1,\ldots,j-1\} \,|\, i_{j'} = i_j \right\} \right) \right)$$

Ways to select one "1" in every column of A such that no two "1"'s share a row

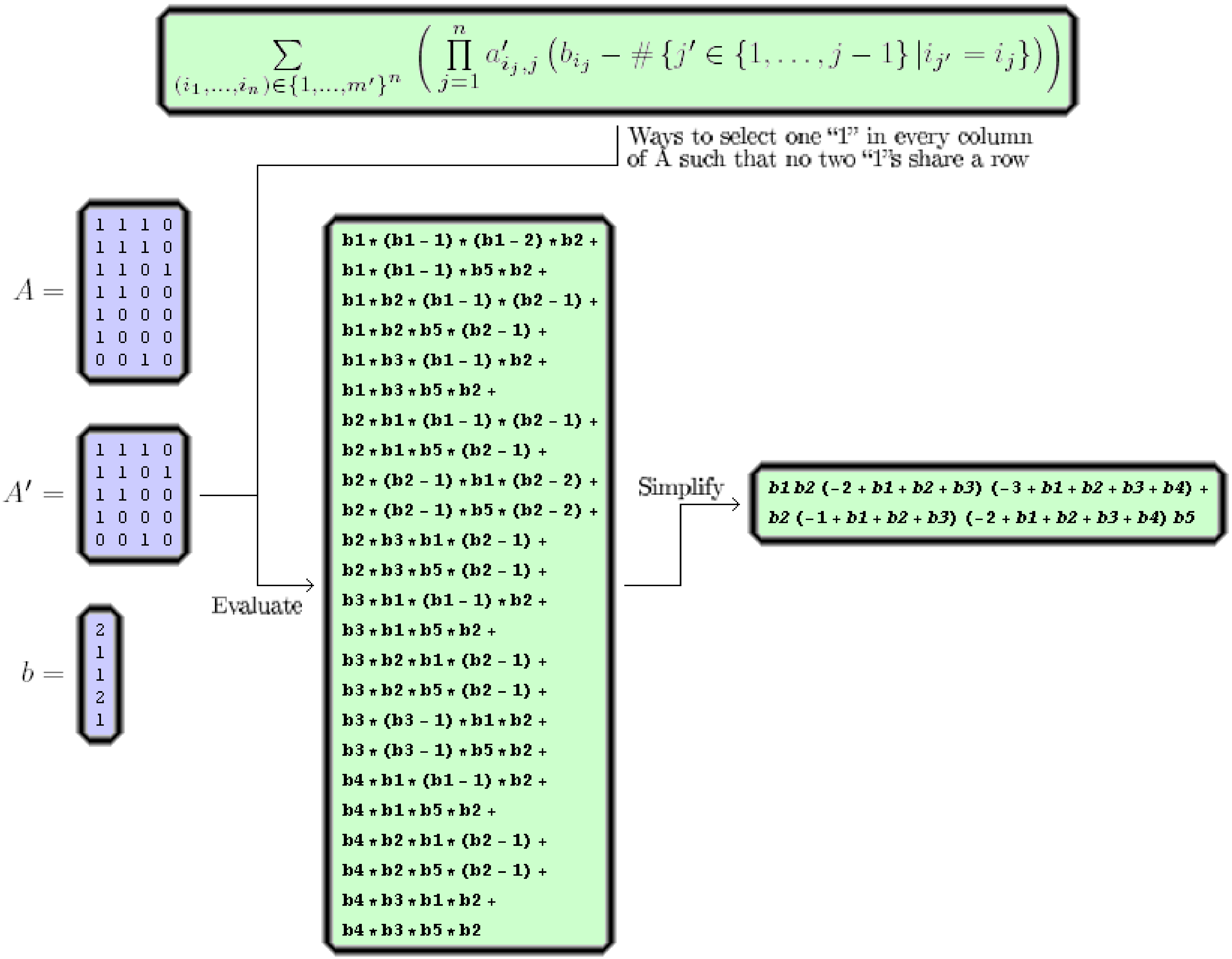
$$A = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

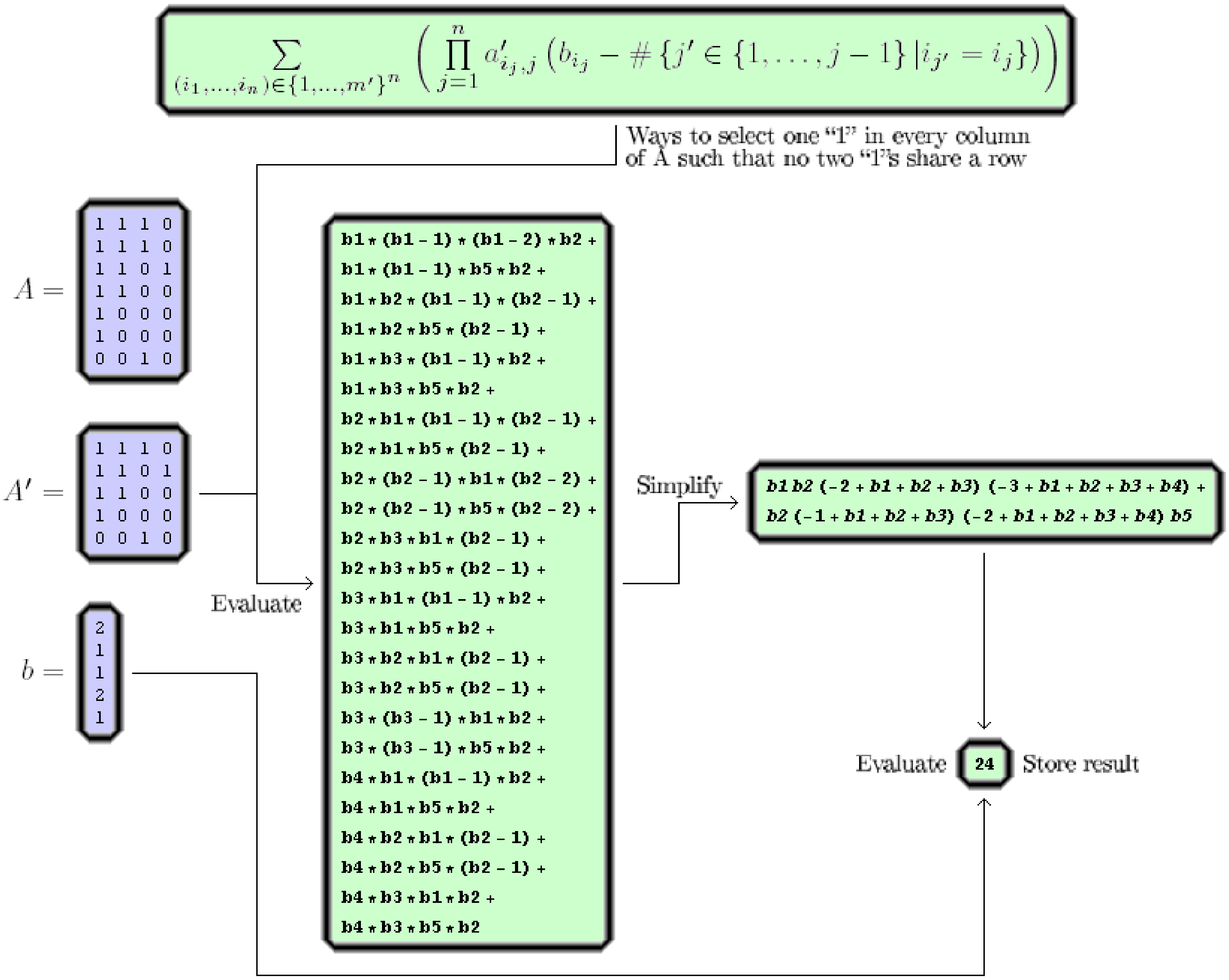
$$A' = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$b = \begin{pmatrix} 2 \\ 1 \\ 1 \\ 2 \\ 1 \end{pmatrix}$$

Evaluate

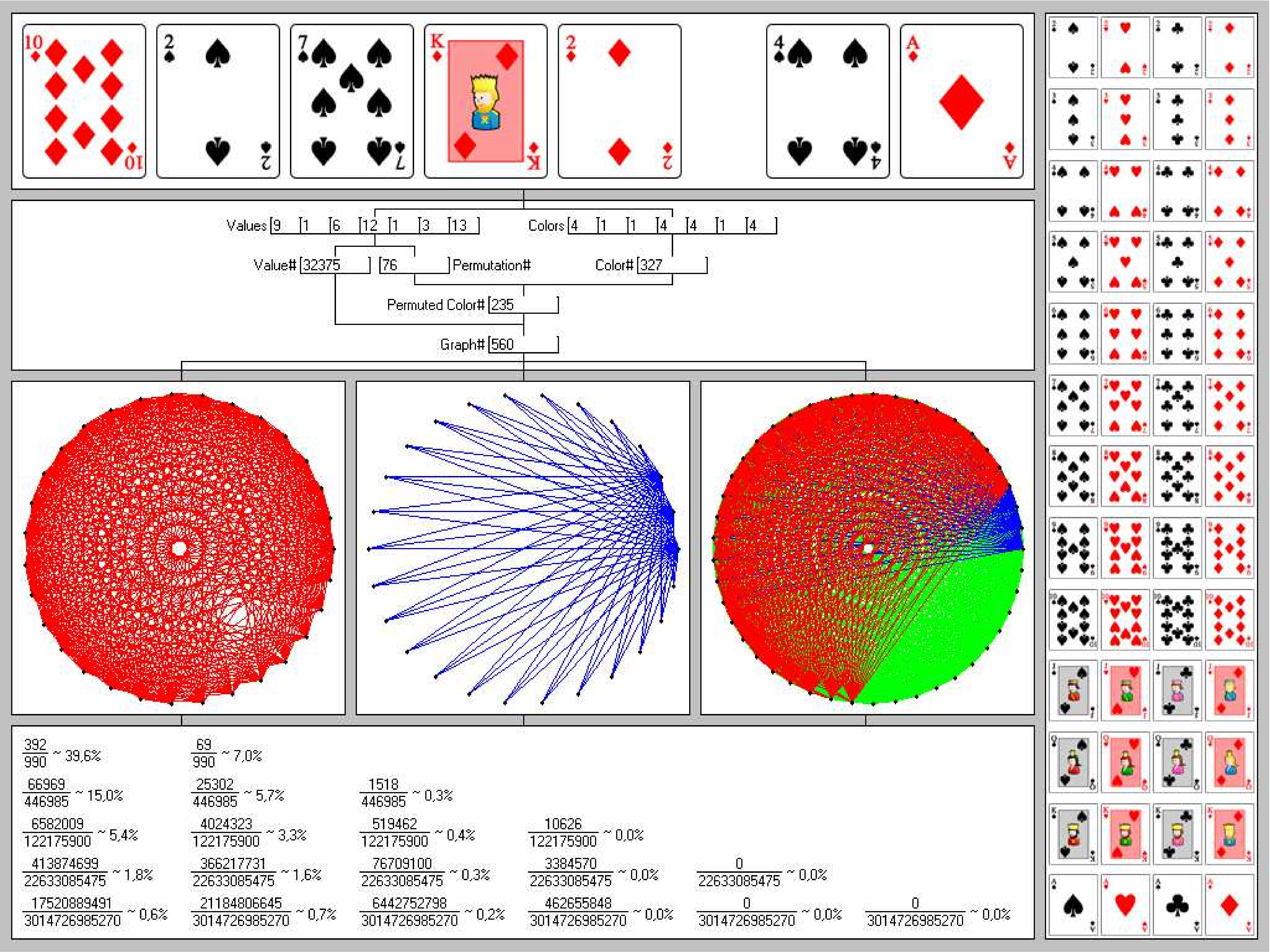
```
b1 * (b1 - 1) * (b1 - 2) * b2 +
b1 * (b1 - 1) * b5 * b2 +
b1 * b2 * (b1 - 1) * (b2 - 1) +
b1 * b2 * b5 * (b2 - 1) +
b1 * b3 * (b1 - 1) * b2 +
b1 * b3 * b5 * b2 +
b2 * b1 * (b1 - 1) * (b2 - 1) +
b2 * b1 * b5 * (b2 - 1) +
b2 * (b2 - 1) * b1 * (b2 - 2) +
b2 * (b2 - 1) * b5 * (b2 - 2) +
b2 * b3 * b1 * (b2 - 1) +
b2 * b3 * b5 * (b2 - 1) +
b3 * b1 * (b1 - 1) * b2 +
b3 * b1 * b5 * b2 +
b3 * b2 * b1 * (b2 - 1) +
b3 * b2 * b5 * (b2 - 1) +
b3 * (b3 - 1) * b1 * b2 +
b3 * (b3 - 1) * b5 * b2 +
b4 * b1 * (b1 - 1) * b2 +
b4 * b1 * b5 * b2 +
b4 * b2 * b1 * (b2 - 1) +
b4 * b2 * b5 * (b2 - 1) +
b4 * b3 * b1 * b2 +
b4 * b3 * b5 * b2
```


$\sum_{(i_1,\ldots,i_n)\in\{1,\ldots,m'\}^n} \left(\prod_{j=1}^{n} a'_{i_j,j} \left(b_{i_j} - \#\{j' \in \{1,\ldots,j-1\} \,|\, i_{j'} = i_j\}\right)\right)$
Ways to select one "1" in every column of A such that no two "1"'s share a row
$A =$
1 1 1 0
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
1 0 0 0
0 0 1 0
$A' =$
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
0 0 1 0
$b =$
2
1
1
2
1
Evaluate
b1 * (b1 - 1) * (b1 - 2) * b2 +
b1 * (b1 - 1) * b5 * b2 +
b1 * b2 * (b1 - 1) * (b2 - 1) +
b1 * b2 * b5 * (b2 - 1) +
b1 * b3 * (b1 - 1) * b2 +
b1 * b3 * b5 * b2 +
b2 * b1 * (b1 - 1) * (b2 - 1) +
b2 * b1 * b5 * (b2 - 1) +
b2 * (b2 - 1) * b1 * (b2 - 2) +
b2 * (b2 - 1) * b5 * (b2 - 2) +
b2 * b3 * b1 * (b2 - 1) +
b2 * b3 * b5 * (b2 - 1) +
b3 * b1 * (b1 - 1) * b2 +
b3 * b1 * b5 * b2 +
b3 * b2 * b1 * (b2 - 1) +
b3 * b2 * b5 * (b2 - 1) +
b3 * (b3 - 1) * b1 * b2 +
b3 * (b3 - 1) * b5 * b2 +
b4 * b1 * (b1 - 1) * b2 +
b4 * b1 * b5 * b2 +
b4 * b2 * b1 * (b2 - 1) +
b4 * b2 * b5 * (b2 - 1) +
b4 * b3 * b1 * b2 +
b4 * b3 * b5 * b2
Simplify
b1 b2 (-2 + b1 + b2 + b3) (-3 + b1 + b2 + b3 + b4) +
b2 (-1 + b1 + b2 + b3) (-2 + b1 + b2 + b3 + b4) b5



$$\sum_{(i_1,\ldots,i_n)\in\{1,\ldots,m'\}^n}\left(\prod_{j=1}^{n} a'_{i_j,j}\left(b_{i_j}-\#\{j'\in\{1,\ldots,j-1\}\,|\,i_{j'}=i_j\}\right)\right)$$
Ways to select one "1" in every column of A such that no two "1"s share a row
A =
1 1 1 0
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
1 0 0 0
0 0 1 0
A′ =
1 1 1 0
1 1 0 1
1 1 0 0
1 0 0 0
0 0 1 0
b =
2
1
1
2
1
Evaluate
b1 * (b1 - 1) * (b1 - 2) * b2 +
b1 * (b1 - 1) * b5 * b2 +
b1 * b2 * (b1 - 1) * (b2 - 1) +
b1 * b2 * b5 * (b2 - 1) +
b1 * b3 * (b1 - 1) * b2 +
b1 * b3 * b5 * b2 +
b2 * b1 * (b1 - 1) * (b2 - 1) +
b2 * b1 * b5 * (b2 - 1) +
b2 * (b2 - 1) * b1 * (b2 - 2) +
b2 * (b2 - 1) * b5 * (b2 - 2) +
b2 * b3 * b1 * (b2 - 1) +
b2 * b3 * b5 * (b2 - 1) +
b3 * b1 * (b1 - 1) * b2 +
b3 * b1 * b5 * b2 +
b3 * b2 * b1 * (b2 - 1) +
b3 * b2 * b5 * (b2 - 1) +
b3 * (b3 - 1) * b1 * b2 +
b3 * (b3 - 1) * b5 * b2 +
b4 * b1 * (b1 - 1) * b2 +
b4 * b1 * b5 * b2 +
b4 * b2 * b1 * (b2 - 1) +
b4 * b2 * b5 * (b2 - 1) +
b4 * b3 * b1 * b2 +
b4 * b3 * b5 * b2
Simplify
b1 b2 (-2 + b1 + b2 + b3) (-3 + b1 + b2 + b3 + b4) +
b2 (-1 + b1 + b2 + b3) (-2 + b1 + b2 + b3 + b4) b5
Evaluate
24
Store result

Step 4: Computation of coefficients


Values 9 1 6 12 1 3 13
Colors 4 1 1 4 4 1 4
Value# 32375
76 Permutation#
Color# 327
Permuted Color# 235
Graph# 560
392/990 ~ 39,6%
69/990 ~ 7,0%
66969/446985 ~ 15,0%
25302/446985 ~ 5,7%
1518/446985 ~ 0,3%
6582009/122175900 ~ 5,4%
4024323/122175900 ~ 3,3%
519462/122175900 ~ 0,4%
10626/122175900 ~ 0,0%
413874699/22633085475 ~ 1,8%
366217731/22633085475 ~ 1,6%
76709100/22633085475 ~ 0,3%
3384570/22633085475 ~ 0,0%
0/22633085475 ~ 0,0%
17520889491/3014726985270 ~ 0,6%
21184806645/3014726985270 ~ 0,7%
6442752798/3014726985270 ~ 0,2%
462655848/3014726985270 ~ 0,0%
0/3014726985270 ~ 0,0%
0/3014726985270 ~ 0,0%