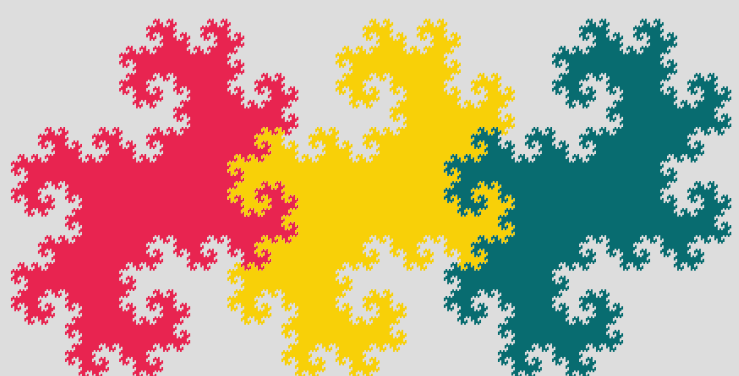




Shift Radix Systems (SRS) with Finiteness Property

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Introduction

For $d \in \mathbb{N}$ and $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{R}^d$ the almost linear mapping

$$\tau_{\mathbf{r}} : \mathbb{Z}^d \mapsto \mathbb{Z}^d$$

$$\mathbf{x} = (x_1, \dots, x_d) \mapsto (x_2, \dots, x_d, -\lfloor \mathbf{r}\mathbf{x} \rfloor)$$

is called the d -dimensional **SRS** associated with $\mathbf{r}$.

$$\mathcal{D}_d := \{\mathbf{r} \in \mathbb{R}^d \mid \text{orbits of } \tau_{\mathbf{r}} \text{ ultimately periodic}\}$$

$$\mathcal{D}_d^{(0)} := \{\mathbf{r} \in \mathbb{R}^d \mid \text{orbits of } \tau_{\mathbf{r}} \text{ end up in } 0\}$$

Elements of $\mathcal{D}_d^{(0)}$ are said to have the **finiteness property**.

Complex analogue (**Gaussian SRS**):
 $\gamma_{\mathbf{r}} : \mathbb{Z}[\mathbf{i}]^d \mapsto \mathbb{Z}[\mathbf{i}]^d$ ($r \in \mathbb{C}$), $\mathcal{G}_d, \mathcal{G}_d^{(0)}$.

Special interest lies in the cases $d = 2$ for SRS and $d = 1$ for GSRS. We identify $\mathbb{C} \leftrightarrow \mathbb{R}^2$ and $\mathbb{Z}[\mathbf{i}] \leftrightarrow \mathbb{Z}^2$ and get

$$\tau_{(r,s)} : \mathbb{Z}^2 \mapsto \mathbb{Z}^2$$

$$(a, b) \mapsto (b, -\lfloor ra + sb \rfloor)$$

$$\gamma_{(r,s)} : \mathbb{Z}^2 \mapsto \mathbb{Z}^2$$

$$(a, b) \mapsto (-\lfloor ra - sb \rfloor, -\lfloor rb + sa \rfloor).$$

Main Goal and Motivation

The main objective is to **characterize** the sets $\mathcal{D}_d^{(0)}$ and $\mathcal{G}_d^{(0)}$. The reason is a strong relation to finiteness properties in other numeration systems.

Definition. For a non-integral, real number $\beta > 1$ the set $\mathcal{A} := \{0, 1, \dots, \lfloor \beta \rfloor\}$ is called the **set of digits** and for every $\gamma \in [0, \infty)$ the unique representation $\gamma = a_m \beta^m + a_{m-1} \beta^{m-1} + \dots$, where $m \in \mathbb{Z}$, $a_i \in \mathcal{A}$, and $0 \leq \gamma - \sum_{i=k}^m a_i \beta^i < \beta^k$ for all $k \leq m$ is called the **greedy expansion** of γ with respect to β . Let $\text{Fin}(\beta)$ denote the set of all $\gamma \in [0, 1)$ having finite greedy expansion with respect to β . β is said to have **property (F)** iff $\text{Fin}(\beta) = \mathbb{Z}[\frac{1}{\beta}] \cap [0, 1)$.

Remark. If β has property (F) then it is an algebraic integer.

Theorem. Let $\beta > 1$ be a non-integral algebraic integer and $(X - \beta)(X^{d-1} + r_{d-2}X^{d-2} + \dots + r_1X + r_0)$ its minimal polynomial. Then β has property (F) $\Leftrightarrow (r_0, \dots, r_{d-2}) \in \mathcal{D}_{d-1}^{(0)}$.

(Akiyama et al. 2005)

Definition. Let $P(X) = X^d + p_{d-1}X^{d-1} + \dots + p_1X + p_0 \in \mathbb{Z}[X]$, $\mathcal{R} := \mathbb{Z}[X]/P(X)\mathbb{Z}[X]$, $\mathcal{N} := \{0, 1, \dots, |p_0| - 1\}$, and $x := X + P(X)\mathbb{Z}[X] \in \mathcal{R}$. $(P, \mathcal{N})$ is called a **CNS**, P a **CNS polynomial** and $\mathcal{N}$ the **set of digits** if every non-zero element $A(x) \in \mathcal{R}$ can be represented uniquely in the form $A(x) = a_mx^m + a_{m-1}x^{m-1} + \dots + a_1x + a_0$.

Theorem. Let $P(X) = X^d + p_{d-1}X^{d-1} + \dots + p_1X + p_0 \in \mathbb{Z}[X]$. Then

$$P \text{ is a CNS polynomial} \Leftrightarrow (\frac{1}{p_0}, \frac{p_{d-1}}{p_0}, \dots, \frac{p_2}{p_0}, \frac{p_1}{p_0}) \in \mathcal{D}_d^{(0)}.$$

(Akiyama et al. 2005)

Main Obstacle

The main problem at characterizing $\mathcal{D}_d^{(0)}$ (and $\mathcal{G}_d^{(0)}$) are so-called **critical points**. It can be shown that $\mathcal{D}_d^{(0)}$ can be gained from $\mathcal{D}_d$ by subtracting certain convex polyhedra. If for a given point infinitely many of these polyhedra are necessary to characterize its neighborhood it is considered critical. In such a case the structure of $\mathcal{D}_d^{(0)}$ tends to be **highly chaotic**.

New Algorithms

The algorithms below **compute** $\mathcal{D}_d^{(0)} \cap H$ ($\mathcal{G}_d^{(0)} \cap H$) where $H \subset \text{int}(\mathcal{D}_d)$ ($\text{int}(\mathcal{G}_d)$) is any convex hull of finitely many points.

Algorithm. (W. [4]) Idea:

- Choose any point $\mathbf{r} \in H$ and compute the corresponding "graph of witnesses" $\mathcal{V}_{\mathbf{r}}$.
- Compute the polyhedron $P_{\mathbf{r}} := \{\mathbf{s} \in \mathbb{R}^d \mid \mathcal{V}_{\mathbf{s}} = \mathcal{V}_{\mathbf{r}}\}$ which is the solution of a system of inequalities.
- Repeat the process with a new $\mathbf{r}$ not contained in any of the previously computed $P_{\mathbf{r}}$ until H is completely covered.
- $\mathcal{D}_d^{(0)} \cap H$ is the union of all $P_{\mathbf{r}}$ for which $\mathbf{r} \in \mathcal{D}_d^{(0)}$.

Properties:

- First "real algorithm" – **terminates for all inputs**.
- **Faster** than previous methods by Brunotte.

Algorithm. (W. [4]) Idea:

- Find a (possibly small) superset V of all sets of witnesses $\mathcal{V}_{\mathbf{r}}$ for $\mathbf{r} \in H$.
- Compute the set of equivalence classes for the equivalence relation: $\mathbf{r} \sim \mathbf{s}$ iff for all $\mathbf{x} \in V$ $\tau_{\mathbf{r}}(\mathbf{x}) = \tau_{\mathbf{s}}(\mathbf{x})$ and $\tau_{\mathbf{r}}^*(\mathbf{x}) = \tau_{\mathbf{s}}^*(\mathbf{x})$.
- Handle classes in sophisticated order to minimize computation time.

Properties:

- **Much faster** than previous methods by Brunotte.

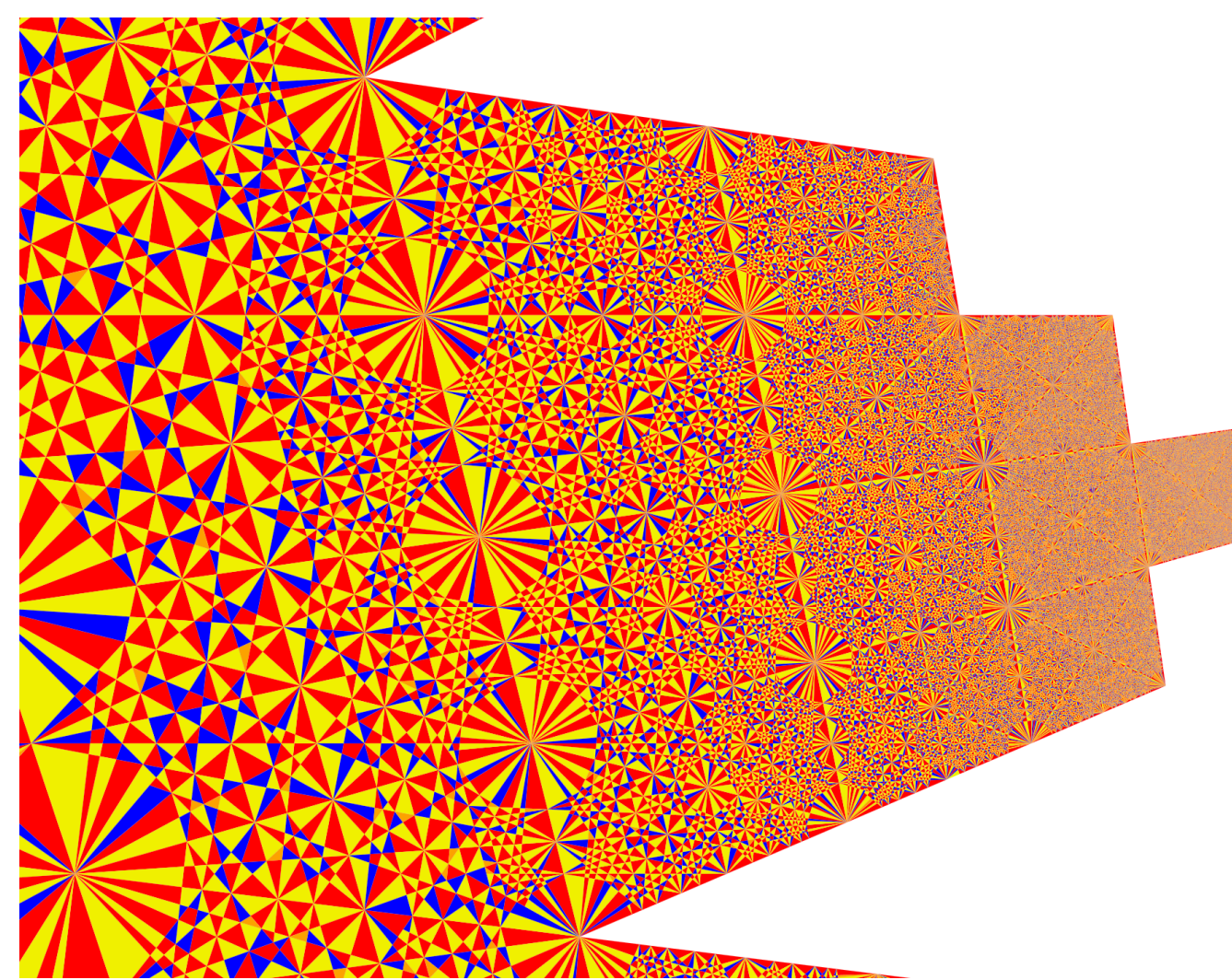


Figure. Hidden structure of $\mathcal{G}_1^{(0)}$ revealed by the equivalence classes of $\sim$.

Standard Methods

For $\mathbf{r} = (r_1, \dots, r_d) \in \mathbb{R}^d$ let $R_{\mathbf{r}}$ denote the companion matrix of $\chi_{\mathbf{r}}(X) = X^d + r_dX^{d-1} + \dots + r_2X + r_1$. Then $\tau_{\mathbf{r}}(\mathbf{x}) = R_{\mathbf{r}}\mathbf{x} + \mathbf{v}_{\mathbf{x}}$ where $\mathbf{v}_{\mathbf{x}} = (0, \dots, 0, \{\mathbf{r}\mathbf{x}\})$. Let $\rho(M)$ denote the spectral radius of a matrix.

- Then
- $\mathcal{D}_d \subseteq \{\mathbf{r} \in \mathbb{R}^d \mid \rho(R_{\mathbf{r}}) \leq 1\}$
 - $\mathcal{E}_d := \{\mathbf{r} \in \mathbb{R}^d \mid \rho(R_{\mathbf{r}}) < 1\} \subseteq \mathcal{D}_d$
 - $\partial\mathcal{D}_d = \{\mathbf{r} \in \mathbb{R}^d \mid \rho(R_{\mathbf{r}}) = 1\}$

For a tuple π of vectors in $\mathbb{Z}^d$ let $\mathcal{P}(\pi)$ denote the set of all $\mathbf{r} \in \mathbb{R}^d$ for which π is a period of $\tau_{\mathbf{r}}$, i.e. $\pi = (\mathbf{x}_1, \dots, \mathbf{x}_n)$, $\tau_{\mathbf{r}}(\mathbf{x}_1) = \mathbf{x}_2, \dots, \tau_{\mathbf{r}}(\mathbf{x}_n) = \mathbf{x}_1$. Then $\mathcal{P}(\pi)$ is a (possibly degenerate) **convex polyhedron** characterized by a finite set of linear inequalities. If $\mathcal{P}(\pi)$ is non-empty it is called a **cutout polyhedron**.

Theorem. $\mathcal{D}_d^{(0)} = \mathcal{D}_d \setminus \bigcup_{\pi \neq 0} \mathcal{P}(\pi)$ (Brunotte 2001)

A set $V \subseteq \mathbb{Z}^d$ is called a **set of witnesses** for $\mathbf{r} \in \mathbb{R}^d$ iff it is stable under $\tau_{\mathbf{r}}$ and $\tau_{\mathbf{r}}^* := -\tau_{\mathbf{r}} \circ (-\text{id}_{\mathbb{Z}^d})$ and contains a generating set of the group $(\mathbb{Z}^d, +)$ which is closed under taking inverses. Every such set of witnesses has the decisive property:

$$\mathbf{r} \in \mathcal{D}_d^{(0)} \Leftrightarrow \forall \mathbf{a} \in V : \exists n \in \mathbb{N} : \tau_{\mathbf{r}}^n(\mathbf{a}) = 0$$

Theorem. For $\mathbf{r} \in \text{int}(\mathcal{D}_d)$ let

$$V_0 := \{(\pm 1, 0, \dots, 0), \dots, (0, \dots, 0, \pm 1)\},$$

$$V_n := V_{n-1} \cup \tau_{\mathbf{r}}(V_{n-1}) \cup \tau_{\mathbf{r}}^*(V_{n-1}) \text{ for } n \in \mathbb{N},$$

$$V_{\mathbf{r}} := \bigcup_{n \in \mathbb{N}_0} V_n.$$

Then $V_{\mathbf{r}}$ is a **finite set of witnesses** for $\mathbf{r}$. (Brunotte 2001)

Results

Theorem.

- $\mathcal{D}_2^{(0)}$ has at least **22 connected components**
- The largest connected component of $\mathcal{D}_2^{(0)}$ has at least **3 holes**.

(W. 2015 [4])

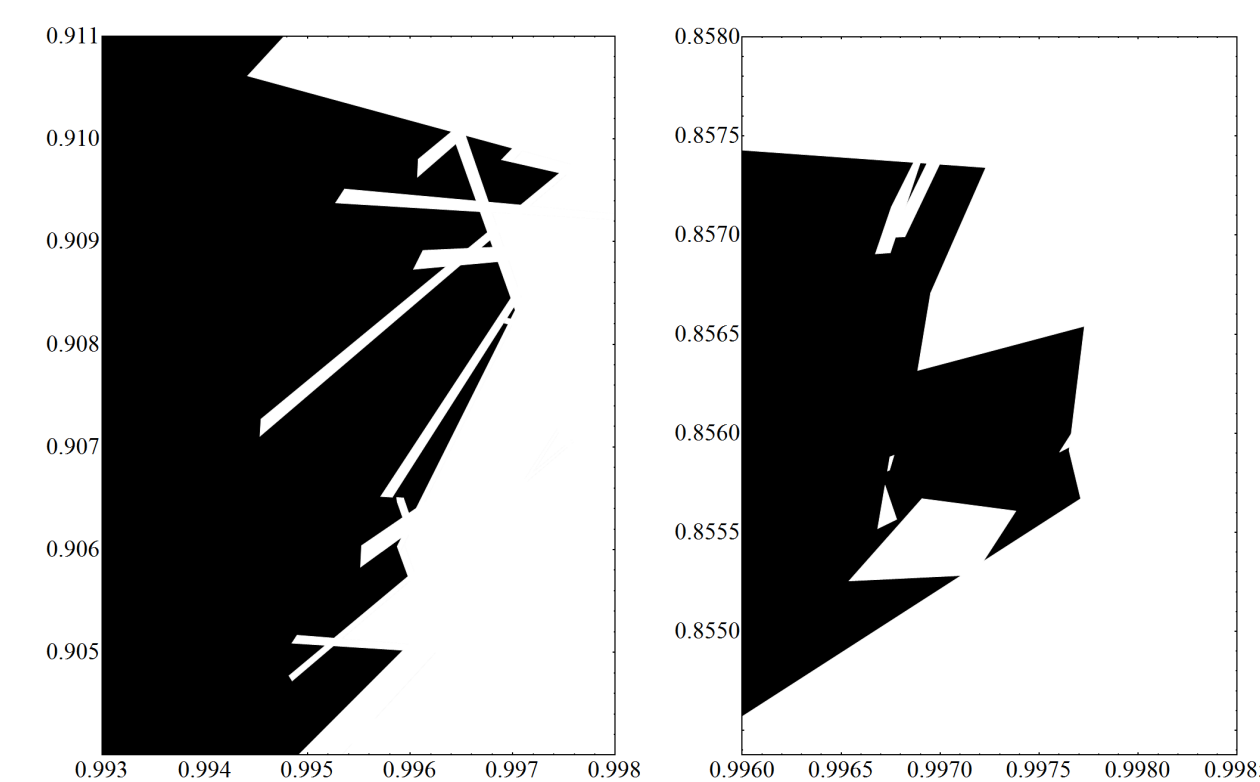


Figure. Four connected components and two holes of $\mathcal{D}_2^{(0)}$.

Conjecture. $\mathcal{G}_1^{(0)} = \mathcal{G}_{\mathbb{C}}$ where $\mathcal{G}_{\mathbb{C}}$ is a (neither open nor closed) polygon given by ten infinite sequences of points in $\mathbb{C}$ (and their complex conjugates). (W.)

Theorem.

- $\mathcal{G}_1^{(0)} \subseteq \mathcal{G}_{\mathbb{C}}$
- $\mathcal{G}_{\mathbb{C}} \cap \{\mathbf{r} \in \mathbb{C} \mid |\mathbf{r}| \leq \frac{1023}{1024}\} \subseteq \mathcal{G}_1^{(0)}$. (W. in preparation [3])

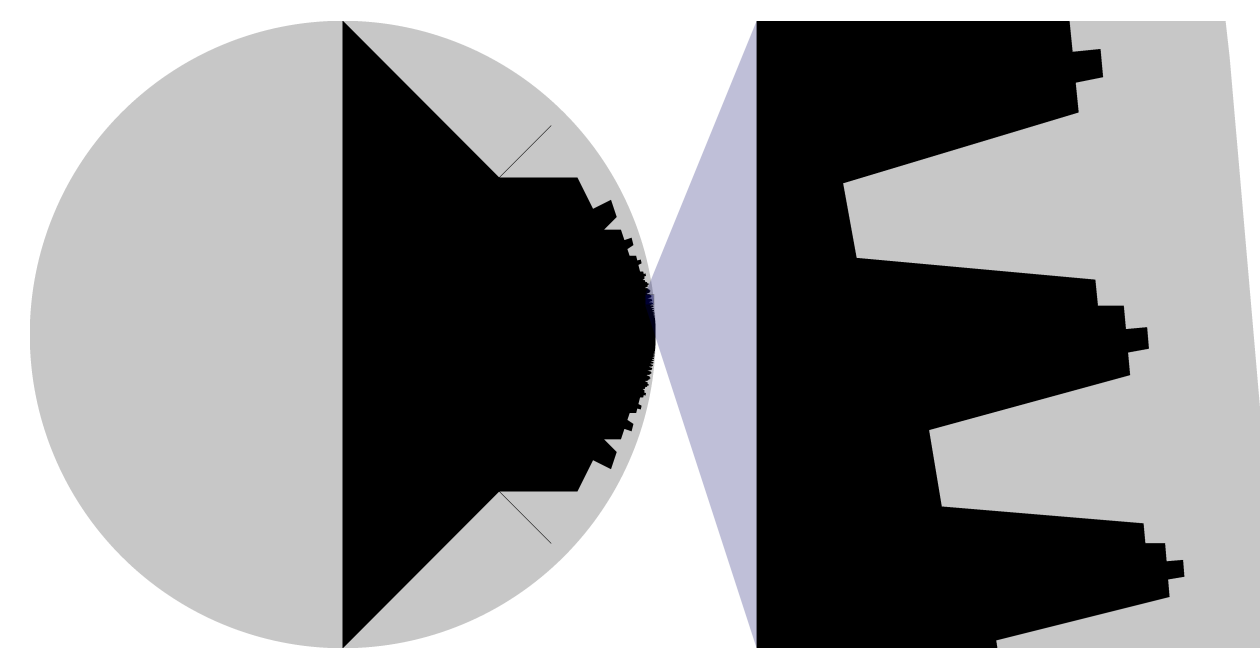


Figure. $\mathcal{G}_1^{(0)}$ in $\mathcal{G}_1$

Joint work (A. Pethő, P. Varga) on **SRS over imaginary quadratic Euclidean domains** in preparation.

Results

Theorem. The quotient $v_d^{(1)}/v_d^{(0)}$ is an integer for all $d \geq 2$. Furthermore we have

$$\frac{v_d^{(1)}}{v_d^{(0)}} = \frac{P_d(3) - 2d - 1}{4} \text{ where}$$

$$P_d(x) := \sum_{k=0}^d \binom{d+k}{2k} \binom{2k}{k} \left(\frac{x-1}{2}\right)^k$$

are the Legendre polynomials. (Kirschenhofer–W. submitted [1])

Contractive Polynomials with given Signature

Definition. Let $\mathcal{E}_d$ denote the set of all coefficient vectors $(a_1, \dots, a_d) \in \mathbb{R}^d$ of contractive polynomials $x^d + a_1x^{d-1} + \dots + a_d \in \mathbb{R}[x]$ (i.e. all roots have absolute value less than 1), and let $\mathcal{E}_d^{(s)}$ denote the subset of the coefficient vectors of those polynomials that have exactly s pairs of complex conjugate roots. Let furthermore $v_d^{(s)} := \lambda_d(\mathcal{E}_d^{(s)})$ denote the d -dimensional Lebesgue measure of $\mathcal{E}_d^{(s)}$.

Conjecture. $v_d^{(s)}/v_d^{(0)} \in \mathbb{Z}$ for all $s \geq 0$ and $d \geq 2s$. (Akiyama–Pethő to appear)

List of Publications

- [1] P. Kirschenhofer and M. Weitzer, *A number theoretic problem on the distribution of polynomials with bounded roots*, submitted.
- [2] A. Pethő, P. Varga, and M. Weitzer, *On shift radix systems over imaginary quadratic Euclidean domains*, in preparation.
- [3] M. Weitzer, *On the characterization of Pethő's loudspeaker*, submitted.
- [4] M. Weitzer, *Characterization algorithms for shift radix systems with finiteness property*, Int. J. Number Theory (2015), to appear.